## Contents

**Introduction/Motivation**  

<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>About this Document</td>
<td>vii</td>
</tr>
<tr>
<td>Purpose(s) of this document</td>
<td>vii</td>
</tr>
<tr>
<td>Text structure</td>
<td>vii</td>
</tr>
<tr>
<td>Compiler messages</td>
<td>viii</td>
</tr>
<tr>
<td>Versions</td>
<td>viii</td>
</tr>
<tr>
<td>Publishing</td>
<td>viii</td>
</tr>
</tbody>
</table>

**1. Teams and Roles**  

<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>Team classes</td>
<td>1</td>
</tr>
<tr>
<td>Role classes and objects</td>
<td>1</td>
</tr>
<tr>
<td>Modifiers for roles</td>
<td>2</td>
</tr>
<tr>
<td>Externalized roles</td>
<td>3</td>
</tr>
<tr>
<td>Protected roles</td>
<td>7</td>
</tr>
<tr>
<td>Type tests and casts</td>
<td>7</td>
</tr>
<tr>
<td>File structure</td>
<td>8</td>
</tr>
<tr>
<td>Acquisition and implicit inheritance of role classes</td>
<td>9</td>
</tr>
<tr>
<td>Acquisition and implicit inheritance of role classes</td>
<td>9</td>
</tr>
<tr>
<td>Regular role inheritance</td>
<td>14</td>
</tr>
<tr>
<td>Name clashes</td>
<td>15</td>
</tr>
<tr>
<td>Team and role nesting</td>
<td>15</td>
</tr>
</tbody>
</table>

**2. Role Binding**  

<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>playedBy relation</td>
<td>19</td>
</tr>
<tr>
<td>Binding interfaces</td>
<td>20</td>
</tr>
<tr>
<td>Legal base classes</td>
<td>20</td>
</tr>
<tr>
<td>Lowering</td>
<td>23</td>
</tr>
<tr>
<td>Lifting</td>
<td>25</td>
</tr>
<tr>
<td>Implicit role creation</td>
<td>26</td>
</tr>
<tr>
<td>Declared lifting</td>
<td>27</td>
</tr>
</tbody>
</table>
5. Team Activation 67
  5.1. Effect of team activation ............................... 67
    5.1.1. Global vs. thread local team activation .......... 67
    5.1.2. Effect on garbage collection ..................... 67
  5.2. Explicit team activation ............................... 67
  5.3. Implicit team activation ............................... 68
  5.4. Guard predicates ...................................... 69
    5.4.1. Regular guards .................................... 70
    5.4.2. Base guards ...................................... 71
    5.4.3. Multiple guards .................................. 73
  5.5. Unanticipated team activation ......................... 74

6. Object Teams API 77
  6.1. Reflection ............................................ 77
  6.2. Other API Elements .................................... 79
  6.3. Annotations .......................................... 81

7. Role Encapsulation 83
  7.1. Opaque roles ......................................... 83
  7.2. Confined roles ...................................... 83

8. Join Point Queries 85
  8.1. Join point queries .................................... 85
  8.2. Query expressions .................................... 85
  8.3. OT/J meta model ..................................... 85

9. Value Dependent Classes 87
  9.1. Defining classes with value parameters .............. 87
  9.2. Using classes with value parameters ................. 87
    9.2.1. Parameter substitution ........................... 87
    9.2.2. Type conformance ................................ 88
  9.3. Restrictions and limitations ......................... 89
Contents

A. OT/J Syntax 91
   A.0. Keywords ........................................... 91
      A.0.1. Scoped keywords .............................. 91
      A.0.2. Inheriting scoped keywords ................. 91
      A.0.3. Internal names .............................. 91
   A.1. Class definitions ................................... 91
   A.2. Modifiers .......................................... 92
   A.3. Method bindings .................................... 92
   A.4. Parameter mappings ............................... 93
   A.5. Statements .......................................... 94
   A.6. Types .............................................. 94
   A.7. Guard predicates .................................... 95
   A.8. Precedence declaration ............................ 95
   A.9. Value dependent types ............................. 95
   A.10. Packages and imports ............................. 96

B. Changes between versions 97
   B.1. Paragraphs changed between versions ............. 97
   B.2. Additions between versions ........................ 99
Introduction

OT/J is an extension to the Java programming language, realizing the concepts of the programming model Object Teams. OT/J introduces aspect-oriented and collaboration/role-based concepts which are smoothly integrated with the object-oriented concepts like inheritance and polymorphism.

Object Teams promotes the notion of collaborations as modules for interacting roles. It does so by introducing two new kinds of modules: teams (§ 1.1.) as higher-order modules for contained roles (§ 1.2.) (see `MyTeam` and its roles `Role1` and `Role2` in Fig. 1).

A `playedBy` relationship (§ 2.1.) binds a role class to a base class, which will be reflected by run-time links between pairs of role and base instances. The main purpose of creating a role-to-base connection is to create a channel for specific communication, which is established by two kinds of method bindings.

A `callin` method binding (§ 4.) intercepts the control flow at a method of the base entity and redirects it to a role method. Looking at Fig. 1, calls to `method2` will be intercepted, causing an invocation of `roleMeth2`. In order to ensure that the effect is purely additive, one of the modifiers `before` or `after` can be used. When specifying a `replace` callin binding, this has approximately the same effect as overriding a method in the context of inheritance. Using the concept of callin bindings, roles unify instance-based inheritance with method call interception as it is used in aspect-oriented programming.

Conversely, a `callout` method binding (§ 3.) simply forwards a method (`roleMeth1`) from a role instance to its `player`, the base instance (`method1`). Callout bindings per se are rather unspectacular, but a callout binding may support `decapsulation`, a term we coined for the inverse of encapsulation. This means that a callout binding can – within limits – access protected or even private members of the corresponding base class (note the “holes” in the border of the base package and its class C2). Decapsulation is also supported for the `playedBy` clause at class level, so that a role may be attached to an otherwise inaccessible class (like a package-visible class or a private inner class).

Several means exist to control when a callin binding is actually effective, i.e., whether or not the base control flow should be intercepted. The most elegant technique is to `activate` or `deactivate` a given team instance (see § 5.), which has the effect that all callin bindings of all contained roles are enabled or disabled in one step. In Fig. 1, activation is symbolized by the switch at the top of `MyTeam`.
About this Document

Levels of this document

Terms, concepts Each chapter of this document starts with a short synopsis of concepts covered by the chapter (like this).
Definition The actual definition is given in small numbered paragraphs.
Examples Examples and accompanying explanations will be interspersed into the definition.

Purpose(s) of this document

This document defines the OT/J programming language. The main goals were to create a precise and complete reference for this language. Didactical considerations had lower priorities, which means that this document is not designed as an introductory tutorial. Still, we advise programmers learning the OT/J language, to consult this document whenever a compiler error message is not perfectly clear to them (see § 0.3).

Text structure

Each chapter of this document starts with a short synopsis of concepts covered by the chapter (see above).

(a) Paragraphs
The actual definition is structured into small paragraphs for easy referral.

(b) Syntax Links
Links to the syntax precede definitions whenever new syntax is introduced.

Interspersed you will find some example program listings. Examples are typeset in a box:

```
public team class MyTeamA {
  
}
```

Explanations for examples look like the following:

Effects:

- Lines 1-3 show a minimal OT/J program, which should not cause any headache.

Examples are given for illustration only.

Additional paragraphs like "Language implementation", or "open issues" provide some background information which is not necessary for understanding the definition, which might however, help to understand why things are the way they are and what other things might be added to the language in the future.
§ 0.3. Compiler messages

Error messages given by the Object Teams compiler refer to this definition whenever appropriate. This way it should be easy to find out, why the compiler rejected your program. Please make sure you are using a language definition whose version matches the version of your compiler.

§ 0.4. Versions

The structure of this document has changed between versions 0.6.1 and 0.7 of this document. This change reflects the transition from our first compiler for OT/J (called otc) and the OTDT (Object Teams Development Tooling) plugin for Eclipse.

Starting with the OTDT v0.7.x, the major and minor number of the tool correspond to the major and minor version number of the OTJLD (this document), ie., the OTDT v1.0.x implements the language as defined in the OTJLD v1.0.

Changes between the current and previous versions are listed in appendix § B.

§ 0.5. Publishing

The sources of this language definition are maintained in a target-independent XML format. Three different versions are generated from these sources, using XSLT:

- Online version\footnote{http://www.objectteams.org/def/} (XHTML)
- Tooling version (XHTML – directly accessible from inside the OTDT)
- Print version (LaTeX/PDF)
Teams and Roles

Fundamental concepts of Teams

Teams and Roles Classes that are defined with the modifier `team` are called team classes, or teams for short. Direct inner classes of a team are called role classes, or roles for short.

Role inheritance Inheritance between teams introduces a special inheritance relationship between their contained roles. The rules of this implicit inheritance are given below (§ 1.3.1.).

Externalized role Roles are generally confined to the context of their enclosing team instance. Subject to specific restrictions, a role may be passed outside its team using the concept of externalized roles (§ 1.2.2.).

Team classes

A class declared with the modifier `team` is a team class (or team for short).

```
public team class MyTeamA {
    ...
}
```

Teams are meant as containers for roles, which are defined in the following paragraphs.

```
public team class MyTeamA {
    public class MyRole
    ...
}
```

Teams introduce a new variant of inheritance for contained role classes (see § 1.3.1., below). Other properties of teams, which are defined in later sections, are:

- Team activation (§ 5.)
- Abstractness and instantiation (§ 2.5.)
- Declared lifting in team methods (§ 2.3.2.)
- Reflective functions defined in `org.objectteams.ITeam` (§ 6.1.)

Apart from these differences, team classes are regular Java classes with methods and fields, whose instances are regular Java objects.

Role classes and objects

Each direct inner class of a team is a role class. Just like inner classes, each instance of a role class has an implicit reference to its enclosing team instance. This reference is immutable.
Within the implementation of a role it can be accessed by qualifying the identifier `this` with the name of the team class, as in:

```
public team class MyTeamA {
    public class MyRole {
        public void print() { System.out.println("Team: " + MyTeamA.this); }
    }
}
```

Creation of role instances is further restricted as defined in § 2.4. Teams can also define role interfaces just like role classes. With respect to role specific properties a role interface is treated like a fully abstract class.

§ 1.2.1. Modifiers for roles

Member classes of a team cannot be static. Also the use of access modifiers for roles is restricted and modifiers have different (stronger) semantics than for regular classes (see below). With respect to accessibility a team acts mainly like a package regarding its roles.

(a) **Role class protection**

A role class must have exactly one of the access modifiers `public` or `protected`. This rule does not affect the class modifiers `abstract`, `final` and `strictfp`.

(b) **protected role classes**

A protected role can only be accessed from within the enclosing team or any of its sub-teams. The actual border of encapsulation is the enclosing team instance. The rules for protected roles are given in § 1.2.3. below.

(c) **public role classes**

Only public roles can ever be accessed outside their enclosing team. Accessing a role outside the enclosing team instance is governed by the rules of **externalized roles**, to be defined next (§ 1.2.2).

(d) **abstract role classes**

A role class has to be marked `abstract` if any of its methods is not effective. The **methods of a role class comprise direct methods and methods acquired by inheritance. In addition to regular inheritance a role class may acquire methods also via implicit inheritance (§ 1.3.1).**

A method may become **effective** by either:

- implementation (i.e., a regular method body), or
- a callout binding (see § 3).

§ 2.5 discusses under which circumstances abstract roles force the enclosing team to be abstract, too.

(e) **Role features**

Access modifiers for members of roles have some special interpretation:
1. A private member is also visible in any implicit sub role (see implicit inheritance § 1.3.1(c)).
In contrast to inner classes in Java, private members of a role are not visible to the enclosing team.

2. The default visibility of role members restricts access to the current class and its sub-classes (explicit and implicit).

3. protected role members can only be accessed from the enclosing team or via callin (§ 4.)

4. public role members grant unrestricted access.

Additionally, a role always has access to all the features that its enclosing team has access to.
Only public members can ever be accessed via an externalized role (§ 1.2.2.)

(f) Static role methods

In contrast to inner classes in pure Java, a role class may indeed define static methods. A static role method requires no role instance but it still requires a team instance in scope. Static role methods can be called:

- from the enclosing team,
- via callin (see § 4.7).

Within a static role method the syntax MyTeam.this is available for accessing the enclosing team instance.

(g) No static initializers

A static field of a role class must not have a non-constant initialization expression. Static initialization blocks are already prohibited for inner classes by Java (see JLS § 8.1.2).

Note: Static initialization generally provides a means for performing initialization code prior to instantiation, i.e., at class-loading time. Before any role can be created already two levels of initialization are performed: (1) The (outer most) enclosing team class performs static initializations when it is loaded. (2) Any enclosing team executes its constructor when it is instantiated. It should be possible to allocate any early initialization to either of these two phases instead of using static role initializers.
§ 1.2.2. or more specifically "family polymorphism", in order to achieve the desired type safety. These theories use special forms of dependent types. Externalized roles have types that depend on a team instance.

§ 1.2.3. deduces even stronger forms of encapsulation from the rules about externalized roles.

(a) Visibility

Only instances of a public role class can ever be externalized.

(b) Declaration with anchored type

Outside a team role types are legal only if denoted relative to an existing team instance (further on called "anchored types"). The syntax is:

```java
final MyTeam myTeam = expression;
RoleClass<@myTeam> role = expression;
```

The syntax `Type<@anchor>` is a special case of a parameterized type, more specifically a value dependent type (§ 9.). The type argument (i.e., the expression after the at-sign) can be a simple name or a path. It must refer to an instance of a team class. The role type is said to be anchored to this team instance.

The type-part of this syntax (in front of the angle brackets) must be the simple name of a role type directly contained in the given team (including roles that are acquired by implicit inheritance).

**Note:**

Previous versions of the OTJLD used a different syntax for anchored types, where the role type was prefixed with the anchor expression, separated by a dot (anchor.Type, see § A.6.). A compiler may still support that path syntax but it should be flagged as being deprecated.

(c) Immutable anchor

Anchoring the type of an externalized role to a team instance requires the team to be referenced by a variable which is marked final (i.e., immutable). The type anchor can be a path `v.f1.f2...` where `v` is any final variable and `f1`... are final fields.

(d) Implicit type anchors

The current team instance can be used as a default anchor for role types:

1. In non-static team level methods role types are by default interpreted as anchored to `this` (referring to the team instance). I.e., the following two declarations express the same:

   ```java
   public RoleX getRoleX (RoleY r) { stmts }
   public RoleX<@this> getRoleX (RoleY<@this> r) { stmts }
   ```

2. In analogy, role methods use the enclosing team instance as the default anchor for any role types.

Note, that `this` and `Outer.this` are always final.

The compiler uses the pseudo identifier `tthis` to denote such implicit type anchors in error messages.
1.2. Role classes and objects

§ 1.2.2. Conformance

Conformance between two types RoleX<@teamA> and RoleY<@teamB> not only requires the role types to be compatible, but also the team instances to be provably the same object. The compiler must be able to statically analyze anchor identity.

§ 1.2.3. Substitutions for type anchors

Only two substitutions are considered for determining team identity:

1. For type checking the application of team methods, this is substituted by the actual call target. For role methods a reference of the form Outer.this is substituted by the enclosing instance of the call target.

2. Assignments from a final identifier to another final identifier are transitively followed, i.e., if t1, t2 are final, after an assignment t1=t2 the types R<@t1> and R<@t2> are considered identical. Otherwise R<@t1> and R<@t2> are incommensurable.

Attaching an actual parameter to a formal parameter in a method call is also considered as an assignment with respect to this rule.

§ 1.2.4. Legal contexts

Anchored types for externalized roles may be used in the following contexts:

1. Declaration of an attribute
2. Declaration of a local variable
3. Declaration of a parameter or result type of a method or constructor
4. In the playedBy clause of a role class (see § 2.1.)

It is not legal to inherit from an anchored type, since this would require membership of the referenced team instance, which can only be achieved by class nesting.

Note:
Item 4. — within the given restriction — admits the case where the same class is a role of one team and the base class for the role of another team. Another form of nesting is defined in § 1.5.

§ 1.2.5. Externalized creation

A role can be created as externalized using either of these equivalent forms:

```java
outer.new Role()
new Role<@outer>()
```

This requires the enclosing instance outer to be declared final. The expression has the type Role<@outer> following the rules of externalized roles. The type Role in this expression must be a simple (unqualified) name.

§ 1.2.6. No import

It is neither useful nor legal to import a role type.

Rationale:
Importing a type allows to use the unqualified name in situations that would
otherwise require to use the fully qualified name, i.e., the type prefixed with its containing package and enclosing class. Roles, however, are contained in a team instance. Outside their team, role types can only be accessed using an anchored type which uses a team instance to qualify the role type. Relative to this team anchor, roles are always denoted using their simple name, which makes importing roles useless.

A static import for a constant declared in a role is, however, legal.

Listing 1: Example code (Externalized Roles):

```java
team class FlightBonus extends Bonus {
  public class Subscriber {
    void clearCredits() { ... }
  }
  void unsubscribe(Subscriber subscr) { ... }
}

class ClearAction extends Action {
  final FlightBonus context;
  Subscriber<@context> subscriber;
  ClearAction (final FlightBonus bonus, Subscriber<@bonus> subscr) {
    context = bonus; // unique assignment to 'context'
    subscriber = subscr;
  }
  void actionPerformed () {
    subscriber.clearCredits();
  }
  protected void finalize () {
    context.unsubscribe(subscriber);
  }
}
```

Effects:

- Lines 1-6 show a terse extract of a published example [NODe02]. Here passengers can be subscribers in a flight bonus program.
- Lines 7-20 show a sub-class of Action which is used to associate the action of resetting a subscriber’s credits to a button or similar element in an application’s GUI.
- Attribute context (line 8) and parameter bonus (line 10) serve as anchor for the type of externalized roles.
- Attribute subscriber (line 9) and parameter subscr (line 10) store a Subscriber role outside the FlightBonus team.
- In order to type-check the assignment in line 12, the compiler has to ensure that the types of LHS and RHS are anchored to the same team instance. This can be verified by checking that both anchors are indeed final and prior to the role assignment a team assignment has taken place (line 11).

Note, that the Java rules for **definite assignments** to final variables ensure that exactly one assignment to a variable occurs prior to its use as type anchor. No further checks are needed.
It is now legal to store this role reference and use it at some later point in time, e.g., for invoking method `clearCredits` (line 15). This method call is also an example for implicit team activation § 5.3.(b).

Line 18 demonstrates how an externalized role can be passed to a team level method. The signature of `unsubscribe` is for this call expanded to

```
void unsubscribe(Subscriber<@context> subscr)
```

(by substituting the call target `context` for `this`). This proves identical types for actual and formal parameters.

### Protected roles § 1.2.3.

Roles can only be public or protected. A protected role is encapsulated by its enclosing team instance. This is enforced by these rules:

(a) **Importing role classes**

   This rule is superseded by § 1.2.2.(i)

(b) **Qualified role types**

   The name of a protected role class may never be used qualified, neither prefixed by its enclosing type nor parameterized by a variable as type anchor (cf. § 1.2.2.(a)).

(c) **Mixing qualified and unqualified types**

   An externalized role type is never compatible to an unqualified role type, except for the substitutions in § 1.2.2.(f) where an explicit anchor can be matched with the implicit anchor `this`.

Rules (a) and (b) ensure that the name of a protected role class cannot be used outside the lexical scope of its enclosing team. Rule (c) ensures that team methods containing unqualified role types in their signature cannot be invoked on a team other than the current team. Accordingly, for role methods the team context must be the enclosing team instance.

(d) **Levels of encapsulation**

   Since protected role types can not be used for externalization, instances of these types are already quite effectively encapsulated by their enclosing team. Based on this concept, encapsulation for protected roles can be made even stricter by the rules of **role confinement**. On the contrary, even protected roles can be externalized as **opaque roles** which still expose (almost) no information. Confinement and opaque roles are subject of § 7.

### Type tests and casts § 1.2.4.

In accordance with § 1.2.2.(e) in OT/J the `instanceof` operator and type casts have extended semantics for roles.
(a) instanceof
For role types the instanceof operator yields true only if both components of the type match: the dynamic role type must be compatible to the given static type, and also type anchors must be the same instance.

(b) Casting
Casts may also fail if the casted expression is anchored to a different team instance than the cast type. Such failure is signaled by a \texttt{org.objectteams.RoleCastException}.

(c) Class literal
A class literal of form \texttt{R.class} is dynamically bound to the class \texttt{R} visible in the current instance context. Using a class literal for a role outside its enclosing team instance (see § 1.2.2.) requires the following syntax:

\begin{verbatim}
RoleClass<@teamAnchor>.class
\end{verbatim}

§ 1.2.5. File structure

Just like regular inner classes, role classes may be inlined in the source code of the enclosing team. As an alternative style it is possible to store role classes in separate role files according to the following rules:

(a) Role directory
In the directory of the team class a new directory is created which has the same name as the team without the .java suffix.

(b) Role files
Role classes are stored in this directory (a). The file names are derived from the role class name extended by .java. A role file must contain exactly one top-level type.

(c) package statement
A role class in a role file declares as its package the fully qualified name of the enclosing team class. The package statement of a role file must use the team modifier as its first token.

(d) Reference to role file
A team should mention in its javadoc comment each role class which is stored externally using a \texttt{@role} tag.

(e) Legal types in role files
The type in a role file must not be an \texttt{enum}.

(f) Imports in role files
A role file may have imports of its own. Within the role definition these imports are visible in addition to all imports of the enclosing team. Only base imports (see § 2.1.2.(d)) must be defined in the team.
Semantically, there is no difference between inlined role classes and those stored in separate role files.

**Note:**
Current Java compilers disallow a type to have the same fully qualified name as a package. However, the JLS does not seem to make a statement in this respect. In OT/J, a package and a type are interpreted as being the same team, if both have the same fully qualified name and both have the `team` modifier.

Listing 2: Role file example:

```java
package org.objectteams.examples;

/*
 * @author Stephan Herrmann
 * @date 20.02.2007
 * @file MyTeamA.java
 */
public team class MyTeamA {
    ...
}
```

```java
package org.objectteams.examples.MyTeamA;

public class MyRole {
    ...
}
```

### Acquisition and implicit inheritance of role classes $\S$ 1.3.

Every team class implicitly implements the predefined interface `org.objectteams.ITeam`. If a team class has no explicit `extends` clause it implicitly extends `org.objectteams.Team`, thus providing implementations for the methods in `org.objectteams.ITeam`. If a team class extends a non-team class, the compiler implicitly adds implementations for all methods declared in `org.objectteams.Team` to the team class. Any subclass of a team (including `org.objectteams.Team`) must again be a team. Interface implementation is not affected by this rule.

Infrastructure provided via interface `org.objectteams.ITeam` is presented in §6.

### Acquisition and implicit inheritance of role classes $\S$ 1.3.1.

A team acquires all roles from its super-team. This relation is similar to inheritance of inner classes, but with a few decisive differences as defined next. Two implementation options are mentioned below (see p. 18), which can be used to realize the special semantics of role acquisition (virtual classes and copy inheritance).
1. TEAMS AND ROLES

Listing 3: Implicit role inheritance

```java
public team class S {
    protected class R0 {...}
    protected class R1 extends R0 {
        boolean ok;
        R2 m() {...}
        void n(R2 r) {...}
    }
    protected class R2 {...}
}

public team class T extends S {
    @Override protected class R1 {
        R2 m() {
            if (ok) { return tsuper.m(); }
            else { return null; }
        }
        void doIt() {
            n(m());
        }
    }
}
```

(a) Role class acquisition

A team T which extends a super-team S has one role class T.R corresponding to each role S.R of the super-team. The new type T.R overrides S.R for the context of T and its roles. Acquisition of role classes can either be direct (see (b) below), or it may involve overriding and implicit inheritance ((c) below).

In the above example (Listing 3) the team S operates on types S.R0, S.R1 and S.R2, while T operates on types T.R0, T.R1 and T.R2.

(Type references like "S.R0" are actually illegal in source code [§ 1.2.3.(b)]. Here they are used for explanatory purposes only)

(b) Direct role acquisition

Within a sub-team T each role S.R of its super-team S is available by the simple name R without further declaration.

The role R2 in Listing 3 can be used in the sub-team T (line 12), because this role type is defined in the super class of the enclosing team.

(c) Overriding and implicit inheritance

If a team contains a role class definition by the same name as a role defined in its super-team, the new role class overrides the corresponding role from the super-team and implicitly inherits all of its features. Such relation is established only by name correspondence.

A role that overrides an inherited role should be marked with an @Override annotation. A compiler should optionally flag a missing @Override annotation with a warning. Conversely, it is an error if a role is marked with an @Override annotation but does not actually override an inherited role.
It is an error to override a role class with an interface or vice versa. A final role cannot be overridden.

Unlike regular inheritance, constructors are also inherited along implicit inheritance, and can be overridden just like normal methods.

In Listing 3 $R_1$ in $T$ implicitly inherits all features of $R_1$ in $S$. This is, because its enclosing team $T$ extends the team $S$ (line 10) and the role definition uses the same name $R_1$ (line 11). Hence the attribute $ok$ is available in the method $m()$ in $T.R_1$ (line 13). $T.R_1$ also overrides $S.R_1$ which is marked by the $\texttt{@Override}$ annotation in line 11.

(d) Lack of subtyping

Direct acquisition of roles from a super-team and implicit inheritance do not establish a subtype relation. A role of a given team is never conform (i.e., substitutable) to any role of any other team. $S.R$ and $T.R$ are always incommensurable.

Note, that this rule is a direct consequence of § 1.2.2.(e).

(e) Dynamic binding of types

Overriding an acquired role by a new role class has the following implication: If an expression or declaration, which is evaluated on behalf of an instance of team $T$ or one of its contained roles, refers to a role $R$, $R$ will always resolve to $T.R$ even if $R$ was introduced in a super-team of $T$ and even if the specific line of code was inherited from a super-team or one of its roles. Only the dynamic type of the enclosing team-instance is used to determine the correct role class (see below for an example).

A special case of dynamically binding role types relates to so-called class literals (see JLS §15.8.2). Role class literals are covered in § 6.1.(c).

The above is strictly needed only for cases involving implicit inheritance. It may, however, help intuition, to also consider the directly acquired role $T.R$ in (b) to override the given role $S.R$.

In line 17 of Listing 3 the implicitly inherited method $n$ is called with the result of an invocation of $m$. Although $n$ was defined in $S$ (thus with argument type $S.R_2$, see line 6) in the context of $T$ it expects an argument of $T.R_2$. This is correctly provided by the invocation of $m$ in the context of $T$.

(f) $tsuper$

$\rightarrow$ Syntax § A.5.4 ($\texttt{TSuperCall}$)

Super calls along implicit inheritance use the new keyword $tsuper$. While super is still available along regular inheritance, a call $tsuper.m()$ selects the version of $m$ of the corresponding role acquired from the super-team.

See § 2.4.2. for $tsuper$ in the context of role constructors.

$tsuper$ can only be used to invoke a corresponding version of the enclosing method or constructor, i.e., an expression $tsuper.m()$ may only occur within the method $m$ with both methods having the same signature (see § 2.3.2.(b) for an exception, where both methods have slightly different signatures).
1.3.1. Teams and roles

In Listing 3, the role \( R_1 \) in team \( T \) overrides the implicitly inherited method \( m() \) from \( S \). \( \text{super} \).\( m() \) calls the overridden method \( m() \) from \( S.R_1 \) (line 13).

(g) Implicitly inheriting super-types

If a role class has an explicit super class (using extends) this relation is inherited along implicit inheritance.

In Listing 3, the role \( R_1 \) in \( T \) has \( T.R_0 \) as its implicitly inherited super class, because the corresponding role in the super-team extends \( R_0 \) (line 3).

Overriding an implicitly inherited super class is governed by § 1.3.2.(b) below. The list of implemented interfaces is merged along implicit inheritance.

(h) Preserving visibility

A role class must provide at least as much access as the implicit super role, or a compile-time error occurs (this is in analogy to JLS §8.4.6.3). Access rights of methods overridden by implicit inheritance follow the same rules as for normal overriding.

(i) Dynamic binding of constructors

When creating a role instance using \( \text{new} \) not only the type to instantiate is bound dynamically (cf. §1.3.1.(e)), but also the constructor to invoke is dynamically bound in accordance to the concrete type.

Within role constructors all \( \text{this}(..) \) and \( \text{super}(..) \) calls are bound statically with respect to explicit inheritance and dynamically with respect to implicit inheritance. This means the target role name is determined statically, but using that name the suitable role type is determined using dynamic binding.

See also §2.5.(a) on using constructors of abstract role classes.

(j) Overriding and compatibility

The rules of JLS §8.4.6 also apply to methods and constructors inherited via implicit inheritance.

(k) Covariant return types

Given a team \( T_1 \) with two roles \( R_1 \) and \( R_2 \) where \( R_2 \) explicitly inherits from \( R_1 \), both roles defining a method \( m \) returning some type \( A \). Given also a sub-team of \( T_1 \), \( T_2 \), where \( T_2.R_1 \) overrides \( m \) with a covariant return type \( B \) (sub-type of \( A \)):

1.3. Acquisition and implicit inheritance of role classes

```java
public team class T1 {
    protected abstract class R1 {
        abstract A m();
    }
    protected class R2 extends R1 {
        A m() { return new A(); }
    }
}
public team class T2 extends T1 {
    protected class R1 {
        @Override B m() { return new B(); } // this declaration renders class T2.R2 illegal
    }
}
```

In this situation role T2.R2 will be illegal unless also overriding `m` with a return type that is at least `B`. Note, that the actual error occurs at the implicitly inherited method T2.R2.`m` which is not visible in the source code, even T2.R2 need not be mentioned explicitly in the source code. A compiler should flag this as an incompatibility at the team level, because a team must specialize inherited roles in a consistent way.

---

### Listing 4: Example code (Teams and Roles):

```java
public team class MyTeamA {
    protected class MyRole {
        String name;
        public MyRole (String n) { name = n; }
        public void print() { System.out.println("id="+name); }
        protected MyRole getRole() { return new MyRole("Joe"); }
    }
    protected MyRole getRole() {
        return new MyRole("Joe");
    }
}
```

```java
public team class MySubTeam extends MyTeamA {
    protected class MyRole {
        int age;
        public void setAge(int a) { age = a; }
        public void print() {
            super.print();
            System.out.println("age="+age);
        }
    }
    public void doit() {
        MyRole r = getRole();
        r.setAge(27);
        r.print();
    }
    ...
    MySubTeam myTeam = new MySubTeam();
    myTeam.doit();
}
```

---

### Listing 5: Program output

```
id=Joe
age=27
```
§ 1.3.2. Regular role inheritance

In addition to implicit inheritance, roles may also inherit using the standard Java keyword `extends`. These restrictions apply:

(a) Super-class restrictions

If the super-class of a role is again a role it must be a direct role of an enclosing team. This rule is simply enforced by disallowing type anchors in the `extends` clause (see § 1.2.2.(g)). As an effect, the super-class may never be more deeply nested than the sub-class.

(b) Inheriting and overriding the `extends` clause

If a role overrides another role by implicit inheritance, it may change the inherited `extends` clause (see § 1.3.1.(g) above) only if the new super-class is a sub-class of the class in the overridden `extends` clause. I.e., an implicit sub-role may specialize the `extends` clause of its implicit super-role.

(c) Constructors and overridden `extends`

Each constructor of a role class that overrides the `extends` clause of its implicit super-role must invoke a constructor of this newly introduced explicit super-class. Thus it may not use a `tsuper` constructor (see § 2.4.2).

(d) Adding implemented interfaces

`implements` declarations are additive, i.e., an implicit sub-role may add more interfaces but has to implement all interfaces of its implicit super-role, too.

(e) Visibility of inherited methods

When a role inherits non-public methods from a regular class (as its super class), these methods are considered as private for the role, i.e., they can only be accessed in an unqualified method call `m()` using the implicit receiver `this`.

Effects:

- According to § 1.3, MyTeamA implements `ITeam` (line 1).
- An implicit role inheritance is created for MySubTeam.MyRole (§ 1.3.1.(c), line 11). If we visualize this special inheritance using a fictitious keyword `overrides` the compiler would see a declaration:

  ```java
  protected class MyRole overrides MyTeamA.MyRole { ... }
  ```

- Invoking `getRole()` on `myTeam` (line 27, 20) creates an instance of `MySubTeam.MyRole` because the acquired role `MyTeamA.MyRole` is overridden by `MySubTeam.MyRole` following the rules of implicit inheritance (cf. § 1.3.1.(e)).
- Overriding of role methods and access to inherited features works as usual.
- As an example for § 1.3.1.(f) see the call `tsuper.print()` (line 15), which selects the implementation of `MyTeamA.MyRole.print`.

Regular role inheritance

In addition to implicit inheritance, roles may also inherit using the standard Java keyword `extends`. These restrictions apply:
1.4. Name clashes

OT/J restricts Java with respect to handling of conflicting names.

(a) Names of role classes

A role class may not have the same name as a method or field of its enclosing team. A role class may not shadow another class that is visible in the scope of the enclosing team.

(b) Names of role methods and fields

Along implicit inheritance, the names of methods or fields may not hide, shadow or obscure any previously visible name.

(see JLS §8.3, §8.4.6.2, §8.5, §9.3, §9.5 (hiding), §6.3.1 (shadowing), §6.3.2 (obscuring).

15

15

Team and role nesting

Multi-level nesting of classes is restricted only by the following rules.

(a) Nested teams

If a role class is also marked using the `team` modifier, it may contain roles at the next level of nesting.

- In the above example (Listing 6) class `RoleAndTeam` starting in line 14 is a role of `OuterTeam` and at the same time a team containing a further role `InnerRole`.

Such a hybrid role-and-team has all properties of both kinds of classes.

(b) Nested classes of roles

A regular role class (i.e., not marked as `team`, see above) may contain local types (see JLS §14.3 - in the example: class `Local`), anonymous types (JLS §15.9.5 - in the example: class defined in lines 18-20) but no member types (JLS §8.5 - in the example: illegal class `IllegalMember`).

The effect is, that nested types of a regular role cannot be used outside the scope of their enclosing role.
(c) **Prohibition of cycles**

A nested team may not extend its own enclosing team.

(d) **Prohibition of name clashes**

A nested team may inherit roles from multiple sources: its explicit super team and any of its implicit super classes (roles) from different levels of nesting. If from different sources a team inherits two or more roles of the same name that are not related by implicit inheritance, this is an illegal name clash.

(e) **Precedence among different supers**

If a role inherits the same feature from several super roles (super and tsuper), an implicitly inherited version always overrides any explicitly inherited feature, i.e., a role with the same simple name is closer related than one with a different name. Also implicit inheritance alone may produce several candidate methods inherited by a role class. This is a result of team-nesting where each level of nesting may add one more tsuper role if outer teams also participate in an inheritance relationship. In this case a role inherited from an *implicit* super team of the enclosing team is closer related than a role inherited from an *explicit* super team. If necessary this rule is applied inside out until a nesting level is found where indeed explicit team inheritance is involved.

So when comparing classes by their fully qualified names the longest common suffix will determine the closest relationship. E.g., `SuperOuter.RoleAndTeamSub.InnerRole` is the closest ancestor of `SubOuter.RoleAndTeamSub.InnerRole` because both share the name suffix `RoleAndTeamSub.InnerRole`.

In the above example (Listing 6) role `OuterTeam.RoleAndTeamSub.InnerRole` has two direct tsuper roles: `OuterTeam.RoleAndTeam.InnerRole` and `SuperOuter.RoleAndTeamSub.InnerRole`. Without the method `foo` defined in lines 27-30, the enclosing class `OuterTeam.RoleAndTeamSub.InnerRole` would inherit the method `foo` defined in `SuperOuter.RoleAndTeamSub.InnerRole` (line 9), because the common name suffix `RoleAndTeamSub.InnerRole` creates a stronger relationship making that class the closest ancestor.
(f) Qualified tsuper

A role in a nested team may qualify the keyword tsuper (see § 1.3.1.(f) above) by a type name in order to select among different implicit super classes. A term OuterTeam.tsuper evaluates to a corresponding implicit super class within the context of the explicit super-class (here: SuperOuter) of the enclosing team "OuterTeam". A method call OuterTeam.tsuper.m() evaluates to the method version within SuperOuter that best corresponds to the current method containing the tsuper-call.

- In the above example [Listing 6] line 28 selects the method version within the superclass of RoleAndTeamSub (i.e., within RoleAndTeam), resolving to OuterTeam.RoleAndTeam.InnerRole.foo().
- Line 29 selects a corresponding method from the context of SuperOuter resolving to SuperOuter.RoleAndTeamSub.InnerRole.foo() which has the same semantics as an unqualified tsuper call would have.

Language implementation:

Role acquisition and implicit inheritance can be implemented in at least two ways.

**Virtual classes**: Each role class is an overridable feature of its enclosing team. Role classes are resolved by dynamic binding with respect to the enclosing team instance. This implementation requires multiple-inheritance in order to also allow regular inheritance between roles of the same team. super and tsuper select parent versions of a method along the two dimensions of inheritance.

**Copy inheritance**: Role acquisition from a super-team has the effect of copying a role definition T.R yielding a new role Tsub.R. All role applications Rx in the role copy refer to Tsub.Rx. Implicit role inheritance extends a role copy in-place. Only the tsuper construct allows to access the previous version of a method (i.e. before in-place overriding).

References:


Role Binding § 2.

Roles and base classes

playedBy relation    A role can be bound to a class outside the team by a playedBy relation, which declares that each role instances is associated to a base instances.

Base class    The class to which a role is bound (using playedBy) is called its base class. Role instances may inherit and override features from their base instance, which is declared using callout (§ 3.) and callin (§ 4.) method bindings.

Bound role    Each role class that declares a playedBy relation is called a bound role. The term bound role may also be used for the instances of such a class.

Lifting / lowering    Translations between a role and its base are called lifting (base to role) § 2.3. and lowering (role to base) § 2.2.

Translation polymorphism    Conformance between a role and a base is governed by translation polymorphism, which refers to a substitutability that is achieved using either lifting or lowering.

Declared lifting    Generally, lifting happens implicitly at data flows between a role object and its base. Team level methods provide additional data flows, where lifting may be declared explicitly.

playedBy relation § 2.1.

(a) Role-base binding

Roles are bound to a base class by the playedBy keyword.

```java
1 public team class MyTeamA {
2    public class MyRole playedBy MyBase {
3        ...
4    }
5 }
```

(b) Inheritance

The playedBy relation is inherited along explicit and implicit (§ 1.3.1.(c)) role inheritance.

(c) Covariant refinement

An explicit sub-role (sub-class using extends) can refine the playedBy relation to a more specific base class (this is the basis for smart lifting (§ 2.3.3.)).

If a role class inherits several playedBy relations from its super-class and its super-interfaces, there must be a most specific base-class among these relations, which is conform to all other base-classes. This most specific base-class is the base-class of the current role.
(d) **No-variance**

An *implicit* sub-role (according to § 1.3.1.(c)) may only add a `playedBy` relation but never change an existing one. Note however, that implicit inheritance may implicitly specialize an existing `playedBy` relation (this advanced situation is illustrated in § 2.7.(d)).

(e) **Use of `playedBy` bindings**

The `playedBy` relation by itself has no effect on the behavior of role and base objects. It is, however, the precondition for translation polymorphism (lowering: § 2.2 and lifting: § 2.3) and for method bindings (callout: § 3 and callin: § 4).

(f) **Effect on garbage collection**

A role and its base object form one conceptual entity. The garbage collector will see a role and its base object as linked in a bidirectional manner. As a result, a role cannot be garbage collected if its base is still reachable and vice versa. Internally a team manages its roles and corresponding bases using weak references. When using one of the `getAllRoles(..)` methods (see § 6.1.(a)), the result may be non-deterministic because these internal structures may hold weak references to objects that will be collected by the next run of the garbage collector. We advise clients of `getAllRoles(..)` to call `System.gc()` prior to calling `getAllRoles(..)` in order to ensure deterministic results.

§ 2.1.1. **Binding interfaces**

Role base bindings may involve classes and/or interfaces. An interface defined as a member of a team is a role interface and may therefore have a `playedBy` clause. Also the type mentioned after the `playedBy` keyword may be an interface.

*Implementation limitation:* The language implementation as of OTDT version 2.0 imposes one particular restriction when binding a role to a base interface: A role binding to a base interface may not contain any callin bindings (§ 4).

§ 2.1.2. **Legal base classes**

Generally, the base class mentioned after `playedBy` must be visible in the enclosing scope (see below (§ 2.1.2.(c)) for an exception). Normally, this scope is defined just by the imports of the enclosing team. For role files (§ 1.2.5.(b)) also additional imports in the role file are considered. § 2.1.2.(d) below defines how imports can be constrained so that certain types can be used as base types, only.

(a) **No role of the same team**

The base class of any role class must not be a role of the same team. It is also not allowed to declare a role class of the same name as a base class bound to this or another role of the enclosing team, if that base class is given with its simple name and resolved using a regular import. Put differently, a base class mentioned after `playedBy` may not be *shadowed* by any role class of the enclosing team.
2.1. playedBy relation

2.1.2. PlayedBy

- **Base imports** as defined below \(\S 2.1.2.(d)\) relax this rule by allowing to import a class as a base class only. In that case no shadowing occurs since the scopes for base classes and roles are disjoint.

**Cycles**

The base class mentioned after playedBy should normally not be an enclosing type (at any depth) of the role class being defined. This rule discourages the creation of cycles where the base instance of a given role \(R\) contains roles of the same type \(R\).

More generally this concerns any sequence of classes \(C_1, C_2, \ldots, C_n\) were each \(C_{i+1}\) is either a member or the base class of \(C_i\), and \(C_n = C_1\). Such structures may be difficult to understand and have certain restrictions regarding callout (\(\S 3.1.(a)\)) and base constructor calls (\(\S 2.4.2.\)). It is furthermore recommended to equip all roles that are played by an enclosing class with a guard predicate (\(\S 5.4.\)) like this:

```java
base when (MyTeam.this == base)
```

This will avoid that the role adapts other instances of the enclosing class which are not the enclosing instance.

It is prohibited to bind a role class to its own inner class.

**Base class decapsulation**

If a base class referenced after playedBy exists but is not visible under normal visibility rules of Java, this restriction may be overridden. This concept is called **decapsulation**, i.e., the opposite of encapsulation (see also \(\S 3.4.\)). A compiler should signal any occurrence of base class decapsulation. If a compiler supports to configure warnings this may be used to let the user choose to (a) ignore base class decapsulation, (b) treat it as a warning or even (c) treat it as an error.

Binding to a \texttt{final} base class is also considered as decapsulation, since a playedBy relationship has powers similar to an extends relationship, which is prohibited by marking a class as \texttt{final}.

Decapsulation is not allowed if the base class is a confined role (see \(\S 7.2.\)). Within the current role a decapsulated base class can be mentioned in the right-hand-side of any method binding (callout (\(\S 3.\)) or callin (\(\S 4.\))). Also arguments in these positions are allowed to mention the decapsulated base class:

- the first argument of one of the role’s constructors (see lifting constructor (\(\S 2.4.1.\)));
- the base side of an argument with declared lifting (see declared lifting (\(\S 2.3.2.\))).

**Base imports**

If the main type in a file denotes a team, the modifier \texttt{base} can be applied to an import in order to specify that this type should be imported for application as a base type only. Example:
import base some_pack MyBase;

public team class MyTeam {
    // simple name resolves to imported class:
    protected class MyRole playedBy MyBase { }
    MyBase illegalDeclaration;  // base import does not apply for this position
}

Types imported by a base import can only be used in the same positions where also base class decapsulation §2.1.2.(c) is applicable. It is recommended that a type mentioned after the keyword playedBy is always imported with the base modifier, otherwise the compiler will give a warning.

Base imports create a scope that is disjoint from the normal scope. Thus, names that are imported as base will never clash with normally visible names (in contrast to §1.4). More specifically, it is not a problem to use a base class’s name also for its role if a base import is used.

(e) No free type parameters

Neither the role class nor the base class in a playedBy binding must have any free type parameters. If both classes are specified with a type parameter of the same name, both parameters are identified and are not considered as free.

From this follows that a role class cannot have more type parameters than its base. Conversely, only one situation exists where a base class can have more type parameters than a role class bound to it: if the role class has no type parameters a generic base class can be bound using the base class’s raw type, i.e., without specifying type arguments.

Note:
The information from the playedBy declaration is used at run-time to associate role instances to base instances. Specifying a base class with free type parameters would imply that only such base instances are decorated by a role whose type is conform to the specified parameterized class. However, type arguments are not available at run-time, thus the run-time environment is not able to decide which base instances should have a role and which should not. This is due to the design of generics in Java which are realized by erasure.

The following example shows how generics can be used in various positions. Note, that some of the concepts used in the example will be explained in later sections.
```java
public class ValueTrafo<T> {
    public T transform(T val) throws Exception { /* ... */ }
}

public team class TransformTeam {
    protected class SafeTrafo<U> playedBy ValueTrafo<U> {
        protected U safeTransform(U val) {
            try {
                return transform(val);
            } catch (Exception e) {
                return val;
            }
        }
    }

    <V> V perform(ValueTrafo<V> as SafeTrafo<V> trafo, V value) {
        return trafo.safeTransform(value);
    }
}

ValueTrafo<String> trafo = new ValueTrafo<String>();
TransformTeam safeTrafo = new TransformTeam();
String s = safeTrafo.perform(trafo, "Testing");
```

**Explanation**

- Line 5 shows a role with type parameter U where the type parameter is identified with the corresponding type parameter of the role’s base class (which is originally declared as T in line 1).
- Line 6 shows a callout binding (§ 3.) which maps a base method to a corresponding role method while maintaining the flexible typing.
- The regular method in lines 7-13 just passes values of type U around.
- The generic method in line 15 ff. uses declared lifting (§ 2.3.2.) to obtain a role for a given base object. The method has no knowledge about the concrete type arguments of either role nor base, but works under the guarantee that both type arguments will be the same for any single invocation.
- Lines 20 ff. finally create instances of base and team and invoke the behavior thereby instantiating type parameters to String.

**Lowering**

Each instance of a bound role class internally stores a reference to its base object. The reference is guaranteed to exist for each bound role instance, and cannot be changed during its lifetime.

(a) **Definition of lowering**

Retrieving the base object from a role object is called lowering. No other means exists for accessing the base reference.

(b) **Places of lowering**

The lowering translation is not meant to be invoked by client code, but implicit translations are inserted by the compiler at all places where a role type is provided while the corresponding base type (or a super type) was expected.
In other words: lowering translations are inserted by the compiler at all places in a
program which would otherwise not be type correct and which using lowering are
statically type correct. This may concern:

- the right hand side of an assignment wrt. the static type of the left hand side,
- the argument values of a method or constructor call wrt. the static type of the
corresponding formal parameter,
- the return value of a method compared to the declared return type of the
method.
- a role parameter in a callout binding (§ 3.3.(d))
- or the return value in a callin binding (§ 4.5.(d))

```
public team class MyTeamA {

    public class MyRole playedBy MyBase {
        . . .
    }

    void useMyBase(MyBase myb) {
        . . .
    }

    MyRole returnMyRole() {
        . . .
    }

    public void doSomething() {
        MyRole r = new MyRole(new MyBase());
        MyBase b = r;
        useMyBase(r);
        MyBase b2 = returnMyRole();
    }
}
```

**Effects:** An instance of type MyRole is lowered to type MyBase when

- assigning it to b (line 7)
- passing it as argument to a method with formal parameter of type MyBase (line
8)
- assigning the return value to a variable of type MyBase (line 9)

*Note:* The constructor call in line 6 uses the *lifting constructor* as defined in § 2.4.1

Lowering translations are not inserted for

- reference comparison (using == or !=)
- instanceof checks
- cast expressions
- return values in callout bindings (§ 3.3.(d))
- parameters in callin bindings (§ 4.5.(d))

For cases where lowering shall be *forced* see § 2.2.(d) below.

(c) Typing

The static type of an implicit lowering translation is the base class declared using
playedBy in the respective role class.

(d) Explicit lowering

If a base type is also the super type of its role, which frequently happens, if a base
reference is known only by the type Object, lowering cannot be deduced automati-
cally, since a type could be interpreted both as a role type and a base type. These
cases may need explicit lowering. For this purpose the role class must declare to implement the interface ILowerable (from org.objectteams.ITeam). This will cause the compiler to generate a method

```java
public Object lower()
```

for the given role class. Client code may use this method to explicitly request the base object of a given role object.

```java
public team class MyTeamA {
    public class MyRole implements ILowerable playedBy MyBase {
        ...
    }
    public void doSomething() {
        MyRole r = new MyRole(new MyBase());
        Object oMyRole = r;
        Object oMyBase = r.lower();
    }
}
```

(e) Lowering of arrays

Lowering also works for arrays of role objects. In order to lower an array of role objects, a new array is created and filled with base objects, one for each role object in the original array. The array may have any number of dimensions at any shape. The lowered array will have exactly the same shape. Note, that each lowering translation will create a new array.

(f) Ambiguous lowering

When assigning a value of a bound role type to a variable or argument of type java.lang.Object this situation is considered as ambiguous lowering because the assignment could apply either (a) a direct upcast to Object or (b) lowering and then upcasting. In such situations the compiler will not insert a lowering translation, but a configurable warning will be issued.

Lifting

Lifting is the reverse translation of lowering. However, lifting is a bit more demanding, since a given base object may have zero to many role objects bound to it. Therefore, the lifting translation requires more context information and may require to create role objects on demand.

(a) Definition of lifting

Retrieving a role for a given base object is called lifting. Lifting is guaranteed to yield the same role object for subsequent calls regarding the same base object, the same team instance and the same role class (see § 2.3.4. for cases of ambiguity that are signaled by compiler warnings and possibly runtime exceptions).

(b) Places of lifting

The lifting translation is not meant to be invoked by client code, but translations are inserted by the compiler at the following locations:

- [Callout bindings (§ 3.3.(c))](result)
Callin bindings (§ 4.5.(a)) (call target and parameters)

Declared lifting (§ 2.3.2.)

(c) Typing

A lifting translation statically expects a specific role class. This expected role class must have a playedBy clause (either directly, or inherited (explicitly or implicitly) from a super role), to which the given base type is conform.

(d) Lifting of arrays

Lifting also works for arrays of role objects. For lifting an array of base objects a new array is created and filled with role objects, one for each base object in the original array. In contrast to the role objects themselves, lifted arrays are never reused for subsequent lifting invocations.

The term translation polymorphism describes the fact that at certain points values can be passed which are not conform to the respective declared type considering only regular inheritance (extends). With translation polymorphism it suffices that a value can be translated using lifting or lowering.

§ 2.3.1. Implicit role creation

Lifting tries to reuse existing role objects so that role state persists across lifting and lowering. If no suitable role instance is found during lifting, a new role is created.

(a) Reuse of existing role objects

A role object is considered suitable for reuse during lifting, if these three items are identical:

1. the given base object
2. the given team object
3. the statically required role type

For the relation between the statically required role type and the actual type of the role object see "smart lifting" (§ 2.3.3.).

(b) Default lifting constructor

Lifting uses a default constructor which takes exactly one argument of the type of the declared base class (after playedBy). By default the compiler generates such a constructor for each bound role. On the other hand, default constructors that take no arguments (as in JLS §8.8.7) are never generated for bound roles.

The super-constructor to be invoked by a default lifting constructor depends on whether the role’s super class is a bound role or not.

- If the super-class is a bound role, the default lifting constructor will invoke the default lifting constructor of the super-class.
- If the super-class is not a bound role, the default lifting constructor will invoke the normal argumentless default constructor of the super-class.

(c) Custom lifting constructor
If a role class declares a custom constructor with the same signature as the default lifting constructor, this constructor is used during lifting. This custom constructor may pre-assume that the role has been setup properly regarding its base-link and registered in the team’s internal map of roles.
If a bound role has an unbound super-class without an argumentless constructor, providing a custom lifting constructor is obligatory, because no legal default lifting constructor can be generated.

(d) Fine-tuning role instantiation
If the lifting operation as defined above degrades the program performance, the lifting semantics can be modified per role class by adding the annotation @org.objectteams.Instantiation which requires an argument of type org.objectteams.InstantiationPolicy in order to select between the following behaviors:

**ONDEMAND** This is the default behavior as defined above.

**ALWAYS** This strategy avoids maintaining the internal role cache, but instead a fresh role instance is created for each lifting request. This may increase the number of role instances but cuts the costs of accessing the cache, which could otherwise become expensive if a cache grows large. As a result of this strategy role state can no longer be shared over time, thus it is discouraged to define fields in a role with this strategy. Also, comparing roles could lead to unexpected results. Therefor, roles with this strategy should implement custom equals and hashCode methods, which should simply delegate to the base instance (using callout § 3).

**NEVER** Roles with this instantiation policy are never instantiated by lifting. Such roles cannot define non-static fields. Otherwise this optimization is fully transparent, specifically callout bindings will refer to the correct base instance.
As of version 2.0 the OT/J compiler does not implement this strategy.

**SINGLETON** Roles declaring this strategy will be instantiated at most once per team. Subsequent lifting requests in the same team will always answer the same role instance. Such roles may receive triggers from callin bindings, but cannot define callout bindings.
As of version 2.0 the OT/J compiler does not implement this strategy.

(3) Declared lifting

(a) Parameters with declared lifting
A non-static team-level method or constructor may declare a parameter with two types in order to explicitly denote a place of lifting. Using the syntax

```java
public void m (BaseClass as RoleClass param) { stmts }
```
a liftable parameter can be declared, provided the second type (RoleClass) is a role of (playedBy) the first type (BaseClass). Furthermore, the role type must be a role of the enclosing team class defining the given method. The role type must be given by its simple (i.e., unqualified) name.
2. ROLE BINDING

Such a signature requires the caller to provide a base object (here BaseClass), but the callee receives a role object (here RoleClass). In fact, the client sees a signature in which the "as RoleClass" part is omitted. Compatibility between caller and callee sides is achieved by an implicitly inserted lifting translation. A signature using declared lifting is only valid, if the requested lifting is possible (see § 2.3.3. and § 2.3.4. for details).

(b) Super in the context of declared lifting

Calling super or tsuper in a method or constructor which declares lifting for one or more parameters refers to a method or constructor with role type parameters, i.e., lifting takes place before super invocation. Nevertheless, the super method may also have a declared lifting signature. It will then see the same role instance(s) as the current method.

(c) Declared lifting of arrays

If a parameter involving explicit lifting should be of an array type, the syntax is

```java
public void m (BaseClass as RoleClass param[]) ...
```

Here the brackets denoting the array apply to both types, BaseClass and RoleClass.

(d) Declared lifting for catch blocks

Also the argument of a catch block may apply declared lifting like in:

```java
catch (BaseException as RoleClass param) { stmts }
```

This syntax is only valid in a non-static scope of a team (directly or nested). In the given example, RoleClass must be played by BaseException. Note, that RoleClass itself need not be a throwable. As the effect of this declaration the catch block will catch any exception of type BaseException and provides it wrapped with a RoleClass instance to the subsequent block. Also note, that re-throwing the given instance param has the semantics of implicitly lowering the role to its base exception before throwing, because the role conforms to the required type Throwable only via lowering.

(e) Generic declared lifting

A method with declared lifting may introduce a type parameter that is bounded relative to a given role type. Such bound is declared as:

```java
<AnyBase base SuperRole>
void teamMethod(AnyBase as SuperRole arg) {
   // body using arg as of type SuperRole
}
```

This means that AnyBase is a type parameter whose instantiations must all be liftable to role SuperRole. The given type bound requires the call site to supply an argument that is compatible to any base class for which the current team contains a bound role that is a sub class of SuperRole, including SuperRole itself. However, Super ROLE itself need not be bound to any base class. On the other hand, different valid substitutions for AnyBase need not be related by inheritance.
Note:
This feature supports generalized treatment of otherwise unrelated base classes.
This is done by defining one bound role for each base under consideration and by having all these roles extend a common unbound role.

Listing 7: Example code (Declared Lifting):

```java
// team class Super {
    public class MyRole playedBy MyBase { ... }
}

// team class Sub extends Super {
    void m (MyRole o) {
        // inside this method o is of type MyRole
        super.m(o);
    }
}

Sub s_team = new Sub();
MyBase b = new MyBase();
// clients see a parameter "MyBase o"
```

Effects:

- Clients use method m with a base instance (type MyBase) as its argument (line 13).
- Before executing the body of m, the argument is lifted such that the method body receives the argument as of type MyRole (line 8).

▷ Smart lifting

In situations where role and base classes are part of some inheritance hierarchies (extends), choosing the appropriate role class during lifting involves the following rules:

(a) Static adjustment

If a base class B shall be lifted to a role class R that is not bound to (playedBy) B, but if a subclass of R — say R2 — is bound to B, lifting is statically setup to use R2, the most general subclass of R that is bound to B or one of its super-types.

Restriction:
This step is not applicable for parameter mappings of replace callin bindings (§ 4.5.(d)).

(b) Dynamic selection of a role class

At runtime also the dynamic type of a base object is considered: Lifting always tries to use a role class that is bound to the exact class of the base object. Lifting considers all role–base pairs bound by playedBy such that the role class is a subclass of the required (statically declared) role type and the base class is a super-class of the dynamic type of the base object.

From those possible pairs the most specific base class is chosen. If multiple role classes are bound to this base class the most specific of these classes is chosen.
(c) Team as closed world

In the above analysis gathering all role-base pairs is performed at compile-time. From this follows, that a team class can only be compiled when all its contained role classes are known and a role class can never be compiled without its team. The analysis includes all roles and their bindings that are inherited from the super-team.

(d) Selection regardless of abstractness

Smart lifting is not affected by abstractness of role classes. For the effect of abstract role classes see § 2.5.

Complex Example:

- If declarations require lifting B3 to R1 this is statically refined to use R2 instead, because this is the most general class declaring a binding to a super–class of B3.
- If the dynamic base type in the same situation is B6, three steps select the appropriate role:
  1. By searching all playedBy clauses (including those that are inherited) the following role–base pairs are candidates: (R2,B2), (R3,B2), (R4,B4) and (R5,B4).
  2. From these pairs the two containing the most specific base class B4 are chosen.
  3. This makes R4 and R5 role candidates, from which the most specific R5 is finally chosen.
If the inheritance hierarchies of the involved base and role classes are given (like in the figure above) the smart lifting algorithm can be rephrased to the following "graphical" rule:

Starting with the dynamic base type \(B_6\) in the example) move upwards the the inheritance relation until you reach a base class bound to a role class indicated by a «playedBy» arrow pointing to the base class \((B_4)\). This role class must be conform to the requested role type.

Switch to the role side along this arrow \((R_4)\). Now move downwards the role inheritance hierarchy as long as the subrole does not refine the playedBy relationship (indicated by another «playedBy» arrow). The bottom role you reach this way \((R_5)\) is the role type selected by smart lifting.

### Binding ambiguities

While all examples so far have only shown 1-to-1 class bindings, several cases of multiple bindings are allowable. Ambiguities may be detected at compile time and/or at runtime.

#### (a) Potential ambiguity

A **potential ambiguity** is given, if two role classes \(R_1\) and \(R_2\) exist such that

- \(R_1\) and \(R_2\) are played by the same base class \(B\), and
- \(R_1\) and \(R_2\) have a common super role \(R_0\), which is also bound to a base class \(B_0\), and
- neither role class \(R_1\) nor \(R_2\) is a (indirect) sub-class of the other.

**Note:** According to § 2.1.(c), if \(B\) is distinct from \(B_0\) it has to be a sub-class of \(B_0\).

**Effect:**

In this case the compiler issues a warning, stating that the \(B\) may not be liftable, because both role classes \(R_1\) and \(R_2\) are candidates and there is no reason to prefer one over the other.

**If no potential ambiguity is detected, lifting will always be unambiguous.**

In the above situation, trying to lift an instance of type \(B\) to the role type \(R_0\) is an **illegal lifting request**. If \(R_0\) is bound to the same base class \(B\) as its sub-roles \(R_1\) and \(R_2\) are, role \(R_0\) is **unliftable**, meaning that no instance of \(R_0\) can ever be obtained by lifting.

#### Listing 8: Example code (Potential Ambiguity):

```java
1  team class MyTeam {
2    public class SuperRole playedBy MyBase { ... }
3    public class SubRoleA extends SuperRole { ... }
4    public class SubRoleB extends SuperRole { ... }
5  }
```

#### (b) Definite ambiguity

A **definite ambiguity** is given if

- the situation of potential ambiguity according to (a) above is given and
lifting is requested (either by method binding or explicitly \[\S 2.3.2\]) from the
shared base class \(B\) to any role class \(R_0\) that is a common super role for \(R_1\) and
\(R_2\).

Definite binding ambiguity also occurs in cases of generic declared lifting\[\S 2.3.2.(e)\]
if the specified role \(R\) is unbound and if two independent sub-roles \(R_1\) and \(R_2\) exist
that introduce a playedBy binding to the same base class \(B_X\). In this case no potential
ambiguity is flagged because roles \(R_1\) and \(R_2\) have no shared bound super-role.

**Effect:**

*Code causing definite ambiguity is required to handle `org.objectteams.LiftingFailedException`.*

In cases of definite binding ambiguity lifting will indeed fail except for some corner
cases. Such corner cases may arise if lifting already finds an appropriate role in the
cache or if an (indirect) subrole of the ambiguously bound role is an unambiguous
lift target for the concrete type of the base object at run-time. See also \[\S 2.3.5\].

**Listing 9: Example code (Definite Ambiguity):**

```java
1  team class MyTeam {
2      public class SuperRole playedBy MyBase { ... }
3      public class SubRoleA extends SuperRole playedBy SubBase { ... }
4      public class SubRoleB extends SuperRole playedBy SubBase { ... }
5          public void useSuperRole(SubBase as SuperRole r) {...} // must declare
6          LiftingFailedException
7      }
8
9  // plus these calls:
10     MyTeam mt = new MyTeam();
11     mt.useSuperRole(new SubBase()); // will throw a LiftingFailedException}
```

(c) Actual ambiguity

At runtime **actual ambiguity** may occur if for the *dynamic type* of a base to be
lifted the conditions of (b) above hold accordingly. Actual ambiguity is only possible
in cases reported by the compiler as potential or definite ambiguity.

**Effect:**

*An actual ambiguity is reported at runtime by throwing a `org.objectteams.LiftingFailedException`.*

**Listing 10: Example code (Actual Ambiguity):**

```java
1  import org.objectteams.LiftingFailedException;
2  team class MyTeam {
3      public class SuperRole playedBy MyBase { ... }
4      public class SubRoleA extends SuperRole playedBy SubBase { ... }
5      public class SubRoleB extends SuperRole playedBy SubBase { ... }
6          public void useSuperRole(MyBase as SuperRole r) throws LiftingFailedException {...}
7  }
8  // plus these calls:
9     MyTeam mt = new MyTeam();
10    mt.useSuperRole(new SubBase()); // will throw a LiftingFailedException
```

(d) Mismatching role

In cases of potential ambiguity another runtime error may occur: a **mismatching role**
is encountered when a role is found in the cache, which is not conform to the
required type. This happens, if the base object has previously been lifted to a type
that is incompatible with the currently requested type.
Effect:
This is reported by throwing a \texttt{org.objectteams.WrongRoleException}.

Listing 11: Example code (Mismatching Role):

```java
import org.objectteams.LiftingFailedException;

team class MyTeam {
  public class SuperRole playedBy MyBase { ... }
  public class SubRoleA extends SuperRole { ... }
  public class SubRoleB extends SuperRole { ... }

  public void useRoleA(MyBase as SubRoleA r) throws LiftingFailedException { ... }
  public void useRoleB(MyBase as SubRoleB r) throws LiftingFailedException { ... }
}

// plus these calls:
MyTeam mt = new MyTeam();
MyBase b = new MyBase();
mt.useRoleA(b); // creates a SubRoleA for b
mt.useRoleB(b); // finds the SubRoleA which is not compatible to the expected type SubRoleB.
```

From the second item of § 2.3.4.(a) follows, that for binding ambiguities different role hierarchies are analyzed in isolation. For this analysis only those role classes are considered that are bound to a base class (directly using \texttt{playedBy} or by inheriting this relation from another role class). I.e., two role classes that have no common bound super role will never cause any ambiguity.

Consequences of lifting problems § 2.3.5.

The rules for lifting and role binding allow (after issuing a warning) two problematic situations:

1. A potential binding ambiguity makes selection of the approprate role type impossible (§ 2.3.4.(a))
2. A role which might be relevant for lifting is abstract (§ 2.5.(b))

Whenever lifting fails for one of these reasons an \texttt{org.objectteams.LiftingFailedException} (§ 6.2.(d)) is thrown. Given that this is a checked exception and depending on the location requiring lifting this has the following consequences:

(a) Problematic declared lifting
A method with declared lifting (§ 2.3.2.) may have to declare \texttt{org.objectteams.LiftingFailedException}.

(b) Problematic callout binding
The role method of a callout binding with result lifting (§ 3.3.(c)) may have to declare \texttt{org.objectteams.LiftingFailedException}.

(c) Problematic callin binding
A callin binding (§ 4.) may silently fail due to a \texttt{org.objectteams.LiftingFailedException}. This exception will actually remain hidden because the callin binding is not explicitly invoked from any source code but implicitly by the runtime dispatch mechanism. To signal this situation the compiler raises an error against such callin binding.
However, the compiler should allow to configure this error and understand the warning token "hidden-lifting-problem" for suppressing this problem [§ 4.1(b)]. If the problem is ignored/suppressed and if at runtime the lifting problem occurs, triggering of the call-in binding will silently fail, i.e., the program will continue in this situation as if the binding hadn’t existed in the first place.

(d) Incompatible redefinition of a role hierarchy

Consider a team T1 with a method m with declared lifting regarding role R, where no lifting problems are detected. Consider next a sub-team T2 which modifies the hierarchy of role R such that lifting to T2.R is problematic due to a binding ambiguity. In this case clients invoking T1.m() could face the situation at runtime that an instance of T2 is used that unexpectedly fails to lift to its role R. Here, the compiler signals a specific error against T2 alerting of the incompatible change.

§ 2.4. Explicit role creation

Lifting is the normal technique by which role objects are created implicitly. This section defines under which conditions a role can also be created explicitly.

§ 2.4.1. Role creation via a lifting constructor

Lifting uses the default constructor for roles (see § 2.3.1). This constructor can be invoked from client code, if the following rules are respected.

(a) Team context

The lifting constructor can be used only within the enclosing team of the role to be instantiated. Thus, qualified allocation expressions (someTeam.new SomeRole(..)) may never use the lifting constructor.

(b) Fresh base object

If the argument to a lifting constructor invocation is a new expression, creating a fresh base object, the use of the lifting constructor is safe. Otherwise the rules of (c) below apply.

(c) Duplicate role runtime check

If it cannot be syntactically derived, that the argument to a lifting constructor is a freshly created base object (b), a compile time warning will signal that an additional runtime check is needed: It must be prevented that a new role is created for a base object, which already has a role of the required type in the given team. It is not possible to replace an existing role by use of the lifting constructor. At runtime, any attempt to do so will cause a org.objectteams.DuplicateRoleException to be thrown. This exception can only occur in situations where the mentioned compile time warning had been issued. § 6.1 will introduce reflective functions which can be used to manually prevent errors like a duplicate role.

§ 2.4.2. Role creation via a regular constructor
Roles may also be created explicitly using a custom constructor with arbitrary signature other than the signature of the lifting constructor. Within role constructors, four kinds of self-calls are possible:

- **base(..)** A constructor of the corresponding base class (§ A.5.(c)), unless the role is involved in base class circularity (§ 2.1.2.(b)), in which case a base constructor call is illegal.

- **this(..)** Another constructor of the same class.

- **super(..)** A constructor of the super-class (normal extends), unless the super-class is bound to a different base class, in which case calling super(..) is not legal.

- **tsuper(..)** A constructor of the corresponding role of the super-team (§ A.5.(e)). Also see the constraint in § 1.3.2.(c).

(a) Unbound roles

Each constructor of a role that is **not bound** to a base class must use one of this(..), super(..) or tsuper(..).

(b) Bound roles

Each constructor of a **bound role** must directly or indirectly invoke either a base(..) constructor or a lifting constructor (see § 2.3.1). Indirect calls to the base constructor or lifting constructor may use any of this(..), super(..) or tsuper(..), which simply delegates the obligation to the called constructor.

If a constructor referenced by base(..) is not visible according to the regular rules of Java, it may still be called using decapsulation (see also § 3.4., § 2.1.2.(c)). Note, that if the super or tsuper role is not bound, delegating the obligation to that unbound role will not work.

(c) Super-call for bound roles

Instead of or prior to calling base(..) a constructor of a bound role explicitly or implicitly calls a super constructor. Which constructor is applicable depends on the super role and its playedBy clause.

- If the super role is bound to the same base class as the current role is,
  - not writing a super-call causes the lifting constructor of the super role to be invoked.
  - explicitly calling a super constructor requires the super constructor to either
    1. create a role instance using a base constructor call (directly or indirectly),
    or
    2. be a lifting constructor receiving a base instance, which the current role must provide as the argument.

- If the super role is bound but the current role refines the playedBy relationship (cf. § 2.1.(c)),
  - a lifting constructor must be called explicitly passing a base object as the argument.

- If the role has an explicit or implicit super role which is unbound the constructor may optionally call a super constructor (using super(..) or tsuper(..)) prior to calling base(..). Otherwise the default constructor is implicitly invoked.
When invoking a lifting constructor of a super role the base object can optionally be obtained by using a base constructor call as an expression:

\[
\text{super(base(<args>))};
\]

The language system evaluates the base constructor by creating an instance of the appropriate base class using a constructor with matching signature. Also the internal links are setup that are needed for accessing the base object from the role and for lifting the base object to the new role in the future.

The syntax for base constructors follows the rule that role implementations never directly refer to any names of base classes or their features.

§ 2.4.3. Role creation in the presence of smart lifting

Explicitly instantiating a role R1 bound to a base B where smart lifting of B to R1 would actually provide a subrole R2 is dangerous: Instantiation enters the R1 into the team’s internal cache. If at any time later lifting this B to R2 is requested, which is a legal request, the runtime system will answer by throwing a org.objectteams.WrongRoleException because it finds the R1 instead of the required R2. For this reason, in this specific situation the explicit instantiation new R1(..) will be flagged by a warning. The problem can be avoided by using R2 in the instantiation expression.

Listing 12: Example code (WrongRoleException):

```java
public class B { void bm() {} }
public team class T {
    protected class R1 playedBy B {...}
    void rm() { /* body omitted */ }
}
public B getDecoratedB () {
    return new R1(new B()); // compile-time warning!
}
public void requestLifting(B as R2 r) {} // plus these calls:
    T t = new T();
    B b = t.getDecoratedB(); // creates an R1 for b
    t.requestLifting(b); // => org.objectteams.WrongRoleException!
```

- A note on line 8: this line passes a fresh instance of B to the lifting constructor of R1 (see § 2.4.1.(b)). In order to return this B instance lowering is implicitly used for the return statement.

- When line 15 is executed, a lifting of b to R2 is requested but due to line 8 an R1 is found in the internal cache.

§ 2.5. Abstract Roles

Overriding of role classes and dynamic binding of role types [§ 1.3.1.(e)] adds new cases to creation with respect to abstract classes.
(a) Using abstract classes for creation

Abstract role classes can indeed be used for object creation. The effect of such a statement is that the team must be marked abstract. Only those sub-teams are concrete that provide concrete versions for all role classes used in creation expressions.

This includes the case, where a super-team has a concrete role class and creates instances of this role class and only the sub-team changes the status of this role class to abstract. Also here the sub-team must be marked abstract, because it contains an abstract role class that is used in creation expressions.

Interpretation:

Since the type in a role creation expression is late-bound relative to the enclosing team instance, abstract role classes can be seen as the hook in a template&hook pattern that is raised from the method level to the class level: A super-team may already refer to the constructor of an abstract role class, only the sub-team will provide the concrete role class to fill the hook with the necessary implementation.

(b) Relevant roles

A team must be marked abstract if one of its relevant roles is abstract.

A role is relevant in this sense if

- the role class is public or if
- an explicit new expression would require to create instances of the role class, or if
- any of the lifting methods of the enclosing team would require to create instances of the role class.

A role is irrelevant with respect to lifting if either of the following holds:

- It is not bound to a base class, neither directly nor by an inherited playedBy clause.
- It has a sub-role without a playedBy clause.
- It is bound to an abstract base class, and for all concrete sub-classes of the base class, a binding to a more specific role class exists.

If neither property, relevance nor irrelevance, can be shown for an abstract role, a warning is given in case the enclosing team is not abstract.

Explicit base references

The role-base link is not meant to be accessed explicitly from programs, but it is fully under the control of compiler and runtime environment. Accessing features of a role’s base object is done by callout bindings (§ 3.). Yet, a keyword base exists, which can be used in the following contexts:

(a) Externalized roles of a base team

If the base class of a role T1.R1 is again a team T2, roles of that team T2 can be externalized (see § 1.2.2) using base as their type anchor. Given that R2 is a role of T2, one could write:
This syntax is only legal within the body of the role $T_1.R_1$ which is bound to the team $T_2$ containing role $R_2$. A static type prefix can be used to disambiguate a base anchor, so the explicit variant of the above type would be $R_2<@R_1.base>$. It is not legal to use a type anchor containing `base` as an element in a path of references like `<@base.field>` or `<@field.base>`.

(b) Explicit base object creation

Within a role constructor (which is not the lifting constructor) the syntax `base(arguments)` causes an instance of the bound base class to be created and linked (see § 2.4.2).

(c) Base call in callin method

Within a callin method (§ 4.2.(d)) an expression `base.m(args)` is used to invoke the originally called method (see § 4.3).

(d) Base guard predicates

Guard predicates (§ 5.4.) can be specified to act on the base side using the `base` when keywords. Within such a base guard predicate `base` is interpreted as a special identifier holding a reference to the base object that is about to be lifted for the sake of a callin method interception (see § 5.4.2.(a)).

(e) Parameter mappings

An expression at the right-hand side of a parameter mapping (parameter in a callin binding (§ 4.4.) or result in a callout binding (§ 3.2.(c)) may use the keyword `base` to refer to the bound base instance. Such usage requires the role method bound in this method binding to be non-static.

(f) Inhibition of modification

In all cases, the base reference is immutable, i.e., `base` can never appear as the left-hand-side of an assignment.

(g) Decapsulation via base reference

In cases § 2.6.(d) and § 2.6.(e) above, members of the base object may be accessed that would not be visible under Java’s visibility rules. Such references are treated as decapsulation in accordance with § 3.4.(a) and § 3.5.(e). Note that accessing a base field via `base` only gives reading access to this field.
Advanced structures

This section discusses how role containment and the playedBy relationship can be combined. It does not define new rules, but illustrates rules defined above. The central idea is that any class can have more than one of the three flavors team, role, and base.

(a) Nesting

If a role (contained in a team) is also a team (marked with the team modifier) it is a nested team. The depth of nesting is not restricted.

(b) Stacking

If the base class to which a role is bound using playedBy is a team, the role is said to be stacked on the base team.

(c) Layering

If roles of a team Secondary are played by roles of another team Primary (i.e., base classes are roles), the team Secondary defines a layer over the team Primary. Such layering requires a final reference anchor from Secondary to an instance of Primary. All playedBy declarations within Secondary specify their base classes anchored to that final link anchor.

Due to the anchored base types, layered teams implicitly support the following guarantee: all base objects of roles of Secondary are contained within the team instance specified by the link anchor. If roles of Secondary contain any callin bindings to non-static base methods, these will be triggered only when a base method is invoked on a base instance contained in the team specified by anchor.

In accordance with § 2.6.(a) the anchor in such anchored playedBy declarations could also be the pseudo identifier base, provided that Secondary is a nested team, which has a playedBy binding to Primary as its base class. This situation is part of the second example below (§ 2.7.(d)) (see T1 playedBy TB1).
(d) Implicit playedBy specialization

According to § 2.1.(d) an implicit sub-role may *implicitly* specialize an existing playedBy relation. This requires the base class to be specified relative to some implicit (OuterTeam.this) or explicit (OuterTeam.base) team anchor. Specializing that team anchor automatically specializes the playedBy declaration, too. This rule never requires any action from a programmer but only explains the interpretation of a playedBy declaration in complex situations.

Two advanced examples demonstrating the above are:

- If a role TOuter1.T.R of a nested team TOuter1.T is played by another role of the outer enclosing team TOuter1.B, subclassing the outer team TOuter1 to TOuter2 will produce a new role TOuter2.T.R which is automatically played by TOuter2.B, an implicit subclass of the original base class TOuter1.B.

- Consider the case where a nested T1 as a role of TOuter is stacked on a base team TB1. Also, T1 is a layered team over TB1 because its role R adapts role TB1.R.

  In this situation the playedBy relation of role TOuter.T1.R is given by a base-anchored type B<@T1.base>. If furthermore TOuter.T1 is subclassed to TOuter.T2 which covariantly refines the inherited playedBy declaration to TB2, then TOuter.T2.R will automatically refine the inherited playedBy relation to TB2.B to follow the new interpretation of the base anchor.
§ 3. Callout Binding

Notion of callout binding

**callout binding** A callout binding declares that a method call to a role object may be **forwarded** to a base method of the associated base object (*the role object *"calls out" to the base*).

**declarative completeness** Even if a role class does not implement all needed methods, but forwards some to its base, also these methods must be declared within the role. Secondly, no forwarding occurs, unless explicitly declared by a callout binding.

**expected/provided** A callout binding binds an **expected** method of the role class (needed but not implemented here) to a **provided** method of the base class.

§ 3.1. Callout method binding

→ [Syntax § A.3.2](#)

A role class may acquire the implementation for any of its (expected) methods by declaring a **callout** binding.

(a) Prerequisite: Class binding

A callout binding requires the enclosing class to be a role class bound to a base class according to § 2.1. However, callout bindings are not allowed if the role is involved in base class circularity (see § 2.1.2.(b)).

(b) Definition

A callout binding maps an abstract role method ("expected method") to a concrete base method ("provided method"). It may appear within the role class at any place where feature declarations are allowed. It is denoted by

```
expected_method_designator -> provided_method_designator;
```

The effect is that any call to the role method will be forwarded to the associated base object using the provided base method.

Listing 13: Example code (Callout):

```java
team class Company {
    public class Employee playedBy Person {
        abstract String getIdentification();
        // callout binding see below...
    }
}
```

(c) Kinds of method designators

A method designator may either be a method name

```
getIdentification -> getName;
```
or a complete method signature including parameter declarations and return type declaration, but excluding any modifiers and declared exceptions.

```
String getIdentification() -> String getName();
```

**Effects:**
- Line 4 declares a callout binding for the role method `getIdentification()`, providing an implementation for the abstract method defined in line 3.
- In combination with the role binding in line 2 this has the following effect:
  - Any call to `Employee.getIdentification` is forwarded to the method `Person.getName`.

Both sides of a callout binding must use the same kind of designators, i.e., designators with and without signature may not be mixed.
Each method designator must uniquely select one method. If a method designator contains a signature this signature must match exactly with the signature of an existing method, i.e., no implicit conversions are applied for this matching. If overloading is involved, signatures must be used to disambiguate.

**(d) Inheritance of role method declarations**
The role method being bound by a callout may be declared in the same class as the binding or it may be inherited from a super class or super interface.

**(e) Callout override**
If an inherited role method is concrete, callout binding regarding this method must use the token "=>" instead of "->" in order to declare that this binding overrides an existing implementation.

Using the "=>" operator for an abstract method is an error.
It is also an error (and not useful anyway) to callout-bind a method that is implemented in the same class as the binding.

**(f) Inheritance of callout bindings**
Callout bindings are inherited along explicit and implicit inheritance. Inherited callout bindings can be overridden using "=>".

**(g) Duplicate bindings**
It is an error if a role class has multiple callout bindings for the same role method.

**(h) Declared exceptions**
It is an error if a base method to be bound by `callout` declares in its `throws` clause any exceptions that are not declared by the corresponding role method.

**(i) Shorthand definition**
A callout binding whose method designators specify full method signatures does not require an existing role method. If no role method is found matching the expected
method of such a callout binding, a new method is implicitly generated. The new
method is static iff the bound base method is static, and it declares the same
exceptions as the bound base method.

A shorthand callout may optionally declare a visibility modifier, otherwise the gener-
ated method inherits the visibility modifier of the bound base method. No further
modifiers are set. If a callout overrides an inherited method or callout, it must not
reduce the visibility of the inherited method/callout.

(j) Inferred callout

If a non-abstract role class inherits an abstract method the compiler tries to infer
a callout binding for implementing the abstract method. Similarly, if a self-call in
a role class cannot be resolved, the compiler tries to infer a callout to resolve the
self-call.

Inference searches for a method in the bound base class such that
1. both methods have the same name
2. both methods have the same number of arguments
3. each argument of the abstract role method is compatible to the corresponding
argument of the base method directly, or using boxing/unboxing or lowering.

Callouts inferred from an interface have public visibility, callouts inferred from a
self-call have private visibility.

Per default inferred callout bindings are disabled, i.e., a compiler must report these as
an error. However, a compiler should allow to configure reporting to produce a warn-
ing only (which can be suppressed using a @SuppressWarnings("inferredcallout")
annotation), or to completely ignore the diagnostic.

(k) Callout to generic method

When referring to a generic base method

\[ \langle T \rangle \ T \ \text{bm}(T \ a) \]

a callout binding may either propagate the method’s genericity as in

\[ \langle T \rangle \ T \ \text{rm}(T \ a) \rightarrow T \ \text{bm}(T \ a) ; \]

or it may supply a valid substitution for the type parameter as in

\[ \text{String rm(String a) \rightarrow String bm(String a)} ; \]

A callout binding either attaches an implementation to a previously declared method or adds
§ 3.1.(i) above) a forwarding method to a role class. Apart from this implementation, callout-
bound methods do not differ from regular methods.

When we say, a callout binding defines forwarding this means that control is passed to the
base object. In contrast, by a delegation semantics control would remain at the role object,
such that self-calls would again be dispatched starting at the role. Callout bindings on their
own do not support delegation. However, in conjunction with method overriding by means of
callin bindings (see § 4.) the effect of delegation can easily be achieved.

▷ Callout parameter mapping

§ 3.2.
§ 3.2. CALLOUT BINDING

(a) with clause
If the method designators in a callout binding are signatures (not just method names), parameters and return value may be mapped by a with(...) sub-clause. Parameter mappings may only occur if the enclosing role is a class, not an interface.

(b) Mapping one parameter
For each parameter of the provided base method, exactly one parameter mapping defines, which value will actually be passed to the base method. Callout parameter mappings have this form:

expression ← base_method_parameter_name

(c) Result mapping
The return value of a callout method may be provided by a result mapping:

result ← expression

The right hand side expression of a result mapping may use the special identifier result to refer to the value returned by the base method.

In a method binding with parameter mappings, it is an error to use result as the name of a regular method argument.

Listing 14: Example code (Callout Parameter Mapping):

```java
1 Integer absoluteValue(Integer integer) ← int abs(int i) with {
  2 integer.intValue() ← i,
  3 result ← new Integer(result)
}
```

(d) Visible names
Each identifier that appears within the expressions of a parameter mapping must be either:

- a feature visible in the scope of the role instance.
- a parameter of the role method (for parameter mappings).
- the special name result (for result mappings).
- in a result mapping also the special name base can be used in order to refer to the bound base instance (provided the method being bound is not static).

The names of base method arguments (i.e., names after mapping) are only legal in the position given in § 3.2.(b).

(e) Implicit parameter mappings
If parameter mappings should be omitted the following conditions must hold:

1. each method parameter of the role method must conform to the corresponding parameter of the base method, and
2. the result type of the base method must conform to the result type of the role method.

Here conformance includes translation polymorphism (cf. § 3.3.(d)).

Parameter correspondence without parameter mapping is determined by declaration order not by names.

Two adjustments can, however, be performed implicitly:

- If the role method has more parameters than the base method, unused trailing parameters may be silently ignored.
- If the role method returns void, any result from the base method may be silently ignored.

Listing 15: Example code (Callout with Parameter Mapping):

```java
public team class MyTeamA {
  public abstract class Role1 {
    abstract void payEuro(float euro);
    abstract float earnEuro();
    void idle(int seconds) { /* do nothing */};
  }
  Role1 boss, worker = // initialization omitted
  public void transaction () {
    boss.payEuro(worker.earnEuro());
    boss.idle(123);
  }
}
public class Staff { // a base class
  public void payDM(float dm) { ... };
  public float earnDM () { ... };
  public int doze () { ... };
  // other methods omitted
}
public team class MySubTeam extends MyTeamA {
  public class Role1 playedBy Staff {
    void payEuro(float euro) -> void payDM(float dm) with {
      euro = 1.95583f -> dm
    }
    float earnEuro() -> float earnDM () with {
      result = result / 1.95583f
    }
    idle => doze; // override existing implementation of idle()
  }
  void doit () {
    transaction();
  }
}
```

Effects:

- Class MyTeamA is declaratively complete and can be type checked because it only uses methods that are visible or declared within this context. MyTeamA.Role1 can, however, not be instantiated, because it is abstract.
- Line 30 has the normal effect of invoking transaction.
• When executing transaction, the call of `worker.earnEuro()` is forwarded to the corresponding base object using method `earnDM()` (binding declaration in line 24). The result is converted by "result / 1.95583f" (line 25).

• Within the same execution of transaction, the call of `boss.payEuro()` is forwarded to the corresponding base object using method `payDM()` (binding declaration in line 21). The parameter `euro` is converted by "euro * 1.95583f" (line 22).

• Method `idle` is forwarded to `doze` without any parameter mapping. This requires `doze` to have a signature that is conformable to the signature of `idle`. In this case a role parameter and a base result are ignored. Using the => operator, this binding overrides the existing implementation of `idle`.

§ 3.3. Lifting and lowering

(For basic definitions see § 2.2. and § 2.3.)

(a) Call target translation
Invoking a base method due to a callout binding first lowers the role object in order to obtain the effective call target.

(b) Parameter translation
Passing a role object as parameter to a callout method implicitly lowers this parameter, if the base method declares a corresponding base type parameter. Lifting of callout parameters is not possible.

(c) Result translation
When returning a base object from a callout method where the role method declares the result to be of a role class, this object is implicitly lifted to the appropriate role. Lowering the result of a callout binding is not possible.

(d) Typing rules
A parameter mapping (implicit by parameter position or explicit by a with clause) is well typed if the left hand side conforms to the right hand side, either by
• type equality
• implicit primitive type conversion
• subtype polymorphism
• translation polymorphism, here: lowering,
• or by a combination of the above.

A result mapping (implicit or explicit by a with clause) is well typed, if the value at the right hand side conforms to the left hand side according to the rules given above, except that translation polymorphism here applies lifting instead of lowering.

(e) Role arrays
For arrays of roles as parameters § 2.2.(e) applies accordingly. For arrays as a return value § 2.3.(d) applies.
Overriding access restrictions

In contrast to normal access restrictions, method bindings may refer to hidden base methods. This concept is the inverse of encapsulation, hence it is called decapsulation. Decapsulation may occur in these positions:

- playedBy declaration (see § 2.1.2.(c))
- base constructor call (see § 2.4.2.(b)).
- callout bindings (see next)
- callout to field (see § 3.5.(e))
- base call within a callin method (see § 4.6.)

(a) Callout to inaccessible base method

By means of callout bindings it is possible to access methods of a base class regardless of their access modifiers. Method bindings are the only place in a program which may mention otherwise inaccessible methods. Access to the callout method at the role side is controlled by regular mechanisms, based on the declaration of the role method.

(b) Sealing against decapsulation

A base package may be "sealed" which re-establishes the standard Java visibility rules. Sealing is achieved by the corresponding capability of Jar files.

(c) Warning levels

A compiler should signal any occurrence of decapsulation. If a compiler supports to configure warnings this may be used to let the user choose to (a) ignore base class decapsulation, (b) treat it as a warning or even (c) treat it as an error (cf. § 2.1.2.(c)). Optionally, a batch compiler may support three levels of verbosity with respect to decapsulation:

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>--nodecapsulation</td>
<td>No warnings.</td>
</tr>
<tr>
<td>default</td>
<td>Warn only if/that access restrictions are overridden.</td>
</tr>
<tr>
<td>--decapsulation</td>
<td>Detailed messages containing the binding and the hidden base method.</td>
</tr>
</tbody>
</table>

(d) Private methods from super classes

If a callout binding shall bind to a private base method, that method must be defined in the exact base class to which the current role class is bound using playedBy. I.e., for private methods § 3.1.(d) does not hold. The same holds for private base fields (see below).

If a private base feature must indeed be callout-bound, a role class must be defined that is played by the exact base class defining the private feature. Another role
bound to a sub-base-class can then be defined as a sub class of the first role. It will inherit the callout binding and through this it can access the desired feature.

1 ```java
   public class SuperBase {
     private int secret;
   }
2 ```

3 ```java
   public class SubBase extends SuperBase {
     /* details omitted */
   }
4 ```

5 ```java
   public team class MyTeam {
     protected class SuperRole playedBy SuperBase {
       int steal() -> get int secret; // OK
     }
6 ```

7 ```java
     protected class SubRole extends SuperRole playedBy SubBase {
       int steal() -> get int secret; // illegal!
     }
9 ```

§ 3.5. Callout to field

Also fields of a base class can be made accessible using a callout binding.

(a) Syntax

Using one of the callout modifiers get or set a role method can be bound to a field of the role’s base class:

1 ```java
   getValue -> get value;
2 ```

3 ```java
   setValue -> set value;
4 ```

where getValue, setValue are abstract role methods of appropriate signatures and value is a field of the bound base class.

A longer syntax is available, too (see line 3 above), which uses complete signatures. For the left hand side § 3.1.(c) applies, for the right hand side, this longer version prepends the field type to the field name.

(b) Compatibility

A role method bound with the modifier get should have no arguments (it may have arbitrary arguments, which are silently ignored) and should have a return type to which the base field is compatible. A role method returning void will ignore the given value and thus has no effect at all, which will be signaled by a compiler warning.

A role method bound with the modifier set must have a first argument that is compatible to the base field’s type (additional arguments - if present - are silently ignored) and must not declare a return type.

(c) Value mapping

Values can be mapped similar to parameter mappings in pure method bindings § 3.2. Such mappings can be used to establish compatibility as required above.

In both get and set bindings, the base side value is denoted by the field’s name (lines 2 and 4 below).

1 ```java
   Integer getValue() -> get int val
   with { result <- new Integer(val) }
3 ```

4 ```java
   void setValue(Integer i) -> set int val
   with { i.intValue() -> val }
5 ```
(d) **Effect**

Callout-binding a role method to a base field generates an implementation for this role method, by which it acts as a getter or setter for the given field of the associated base object.

(e) **Access control**

For accessing an otherwise invisible field, the rules for decapsulation (§ 3.4.) apply accordingly.

Recall, that according to JLS § 8.3 fields may be hidden in sub-classes of a given base class. Therefore, it is relevant to know that a callout to a field will always access the field that is visible in the exact base class of the role class defining the callout. This is especially relevant for accessing private fields.

(f) **Shorthand definition**

Just as in § 3.1.(i) a shorthand definition allows to introduce a callout field access method without prior abstract declaration. This requires the callout field binding to specify types as in line 3 of § 3.5.(a) above. The generated access method is static iff the bound base field is static.

A shorthand callout to field may optionally declare a visibility modifier, otherwise the generated method inherits the visibility modifier of the bound base field. No further modifiers are set. If a callout to field overrides an inherited method or callout, it must not reduce the visibility of the inherited method/callout.

(g) **Callout override**

Similar to method callouts a callout to field may override an existing role method if and only if the token `=>` is used instead of `->` (see § 3.1.(e) and § 3.1.(f)).

(h) **Inferred callout**

If a statement or expression within the body of a bound role class uses a simple name or a name qualified by `this` which can not be resolved using normal rules, the compiler may infer to use a callout to field instead, given that a field of the required name can be found in the role’s declared baseclass.

If a callout to field has explicitly been declared it is used for the otherwise unresolved name, if and only if:

- the callout declares a role method name the is constructed from the token "set" for a setter or "get" for a getter plus the field name with capital first letter,
- the base field referenced by the callout has exactly the required name, and
- the callout kind (set/get) matches the application of the unresolved name as either the left-hand side of an assignment (set) or as an expression (get).

If a matching callout to field is not found, the compiler generates one automatically, which has private visibility.

If a callout to field has been inferred it is an error to directly invoke the implicitly generated callout accessor that is formed according to the above rules.

---

Per default inferred callout bindings are disabled, i.e., a compiler must report these as an error. However, a compiler should allow to configure reporting to produce a warning only (which can be suppressed using a `@SuppressWarnings("inferredcallout")` annotation), or to completely ignore the diagnostic. See also § 3.1.(j).
Callin Binding

Notion of callin binding

Callin bindings realize a forwarding in the direction opposite to callout bindings (see \[3\]). Both terms are chosen from the perspective of a role, which controls its communication with an associated base object. Technically, callin bindings are equivalent to weaving additional code (triggers) into existing base methods.

**Callin** Methods of a base class may be intercepted by a callin binding (the base method "calls into" the role).

**Before/after/replace** The modifiers before, after, replace control the composition of original method and callin method.

**Activation** Callin bindings may be active or inactive according to \[5\].

Callin method binding

(a) Method call interception

A role method may intercept calls to a base method by a callin binding.

(b) Prerequisite: Class binding

A callin binding requires the enclosing class to be a role class bound to a base class according to \[2.1\]. An unliftable role (see \[2.3.4.(a)\]) should not define callin bindings. In this case callin bindings can only safely be introduced in sub-roles which (by an appropriately refined playedBy clause) disambiguate the lifting translation. For corner cases the above rule can be overridden by suppressing the corresponding error using the "hidden-lifting-problem" token (see \[2.3.5\]). This will allow callin bindings to be defined even for unliftable roles expecting that lifting may still succeed by one of the patterns described in \[2.3.4.(b)\].

(c) Callin declaration

A callin binding composes an existing role method with a given base method. It may appear within the role class at any place where feature declarations are allowed. It is denoted by

```
role_method_designator <- callin_modifier base_method_designator;
```

Just like with callout bindings, method designators may or may not contain parameters lists and return type but no modifiers; also, each method designator must exactly and uniquely select one method (cf. \[3.1.(c)\]).

For callin modifiers see below (\[4.2\]).
(d) Multiple base methods
Base method designators may furthermore enumerate a list of methods. If multiple base methods are bound in one callin declaration generally all signatures in this binding must be conform.
However, extraneous parameters from base methods may be ignored at the role.
For result types different rules exist, depending on the applied callin modifier (see next).

(e) Named callin binding
Any callin binding may be labeled with a name. The name of a callin binding is used for declaring precedence (§ 4.8.). A named callin binding overrides any inherited callin binding (explicit and implicit § 1.3.1) with the same name.
It is an error to use the same callin name more than once within the same role class.

(f) Callin to final
When binding to a final base method, the enclosing role must be played by the exact base class declaring the final method. I.e., callin binding to a final method inherited from the base class’s super-class is not allowed. This is motivated by the fact that no sub-class may have a version of a final method with different semantics.

(g) Declared exceptions
It is an error if a role method to be bound by callin declares in its throws clause any exceptions that are not declared by the corresponding base method(s).

(h) Method of enclosing class
In a before or after callin binding the left hand side may alternatively resolve to a method of an enclosing class rather than the current role.

§ 4.2. Callin modifiers (before, after, replace)

(a) Method composition
The kind of method composition is controlled by adding one of the modifiers before, after or replace after the "<-" token of the binding declaration.

(b) Additive composition
The before and after modifiers have the effect of adding a call to the role method at the beginning or end of the base method, resp.
In this case no data are transferred from the role to the base, so if the role method has a result, this will always be ignored.

Listing 16: Example code (Callin):

```java
team class Company {
  protected class Employee playedBy Person {
    public void recalculateIncome() { ... }
    recalculateIncome <- after haveBirthday; // callin binding
  }
}
```
Line 4 declares a callin binding for the role method `recalculateIncome()` defined in line 3. In combination with the role binding in line 2 this has the following effect:

- **After** every call of the method `Person.haveBirthday` the method `Company.recalculateIncome` is called.

(c) Replacing composition

The `replace` modifier causes only the role method to be invoked, replacing the base method.

In this case, if the base method declares a result, this should be provided by the role method. Special cases of return values in callin bindings are discussed in § 4.3.(e)

(d) Callin methods

Role methods to be bound by a callin replacement binding must have the modifier `callin`. This modifier is only allowed for methods of a role class.

A method with the `callin` modifier can only be called

- via a callin replace binding
- by a super or `tsuper` call from an overriding callin method.

It is illegal for a callin method

- to be called directly,
- to be bound using a callout binding, and
- to be bound to a base method using a `before` or `after` callin binding.

Despite these rules a second level role — which is played by the current role — can intercept the execution of a callin method using any form of callin binding.

A callin method cannot override a regular method and vice versa, however, overriding one callin method with another callin method is legal and dynamic binding applies to callin method just like regular methods.

A callin method must not declare its visibility using any of the modifiers `public`, `protected` or `private`. Since callin methods can only be invoked via callin bindings such visibility control would not be useful.

Base calls

Role methods with a `callin` modifier should contain a `base call` which uses the special name `base` in order to invoke the original base method (original means: before replacement).

(a) Syntax

The syntax for base calls is `base.m()`, which is in analogy to super calls. A `base.m()` call must use the same name and signature as the enclosing method. This again follows the rule, that roles should never explicitly use base names, except in binding declarations.
(b) Missing base call
For each callin method, the compiler uses some flow analysis to check whether a base call will be invoked on each path of execution (analysis is very similar to the analysis for definite assignment regarding final variables - JLS §16). The compiler will issue a warning if a base call is missing either on each path (definitely missing) or on some paths (potentially missing). Instead of directly invoking a base call, a callin method may also call its explicit or implicit super version using super.m() or tsuper.m() (see §1.3.1(f)). In this case the flow analysis will transitively include the called super/tsuper version.

(c) Duplicate base call
If a callin method contains several base calls, the compiler gives a warning if this will result in duplicate base call invocations on all paths (definitely duplicate) or on some paths (potentially duplicate). Again super/tsuper calls are included in the flow analysis (see 4.3(b)).

(d) Parameter tunneling
If a base method has more parameters than a callin method to which it is composed, additional parameters are implicitly passed unchanged from the original call to the base call (original means: before interception). I.e., a call base.m() may invisibly pass additional parameters that were provided by the caller, but are hidden from the role method.

(e) Fragile callin binding
If a role method returns void, but the bound base method declares a non-void result, this is reported as a fragile callin binding: The result can still be provided by the base call, but omitting the base call may cause problems depending on the return type:

- For reference return types null will be returned in this case.
- In the case of primitive return types this will cause a ResultNotProvidedException at run-time.

It is an error if a callin method involved in a fragile callin binding has definitely no base call.

(f) Base super calls
If a callin method rm is bound to a base method B1.m that in turn overrides an inherited method B0.m (B0 is a super class of B1), the callin method may use a special form of a base call denoted as

    base.super.rm();

Such base super call invokes the super method of the bound base method, here B0.m. This invocation is not affected by any further callin binding.

A base super call bypasses both the original method B1.m and also other callin bindings that would be triggered by a regular base call. For this reason any application of this construct is flagged by a decapsulation warning (see §3.4).

Comment:
Base calls can occur in callin methods that are not yet bound. These methods have no idea of the names of base methods that a sub-role will bind to them. Also multiple base methods may be bound to the same callin method. Hence the use of the role method’s own name and signature. The language implementation translates the method name and signature back to the base method that has originally been invoked.

Listing 17: Example code (Base Call):

```java
public class ValidatorRole playedBy Point {
    callin void checkCoordinate(int value) {
        if (value < 0)
            base.checkCoordinate(-value);
        else
            base.checkCoordinate(value);
    }
    checkCoordinate <- replace setX, setY;
}
```

Effects:
- Line 2 defines a callin method which is bound to two methods of the base class Point (see line 8).
- The value passed to either setX or setY is checked if it is positive (line 3).
- Lines 4 and 6 show calls of the original method (base calls). While line 6 passes the original value, in the negative case (line 4) the passed value is made positive.

Callin parameter mapping § 4.4.

(a) General case parameter mapping
The rules for mapping callin parameters and result type are mainly the same as for callout bindings § 3.2. except for reversing the -> and <- tokens and swapping left hand side and right hand side.
Callin bindings using before have no result mapping. For result in after callin bindings see § 4.4.(c) below.

(b) Restrictions for callin replace bindings
The right-hand side of a parameter mapping may either be the simple name of a base method argument without further computation, or an arbitrary expression not containing any base method argument.
Each base method argument must either appear as a simple name in exactly one parameter mapping or not be mapped at all. In the latter case, the original argument is "tunneled" to the base call, meaning, the callin method does not see the argument, but it is passed to the base method as expected.
If the base method declares a result, then
- if the role method also declares a result, result must be mapped to itself:
  result -> result
• if the role method does not declare a result, an arbitrary expression may be mapped to result:
  
  \[
  \text{expression} \rightarrow \text{result}
  \]

  If in this situation no result mapping exists, the result of the base call is "tunneled" and passed to the original caller (see fragile callin binding (§ 4.3.(e)) above).

  These rules ensure that these bindings are reversible for the sake of base calls (§ 4.3.). As stated above a fragile callin binding (§ 4.3.(e)) is not allowed with a callin method that definitely has no base call (§ 4.3.(b)). A callin replace binding is not fragile if it provides the base result using a result mapping.

  A callin method bound with replace to a base method returning void must not declare a non-void result.

(c) Mapping the result of a base method

In an after callin binding, the right-hand side of a parameter mapping may use the identifier result to refer to the result of the base method. An after callin binding can, however, not influence the result of the base method, thus mappings with the \( \rightarrow \) token are not allowed for after callin bindings. For before mappings using the \( \rightarrow \) token is already ruled out by § 4.4(a).

(d) Multiple base methods

A callin binding listing more than one base method may use parameter mappings with only the following restriction: if any base parameter should be mapped this parameter must have the same name and type in all listed base method designators. However, different parameter mappings for different base methods bound to the same role method can be defined if separate callin bindings are used.

§ 4.5. Lifting and lowering

For basic definition see § 2.2. and § 2.3. (The following rules are reverse forms of those from § 3.3.)

(a) Call target translation

Invoking a role method due to a callin binding first lifts the base object to the role class of the callin binding, in order to obtain the effective call target. This is why callin bindings cannot be defined in roles that are unliftable due to potential binding ambiguity (see § 4.1(b) above and § 2.3.4(a)).

(b) Parameter translation

During callin execution, each parameter for which the role method expects a role object is implicitly lifted to the declared role class.

(c) Result translation

Returning a role object from a callin method implicitly lowers this object.

(d) Typing rules

A parameter mapping (implicit by parameter position or explicit by a with clause) is well typed if the right hand side conforms to the left hand side, either by
• type equality
• implicit primitive type conversion
• subtype polymorphism
• translation polymorphism, here: lifting;
  however, within replace bindings step 1 of the smart lifting algorithm (§2.3.3.(a))
  is not applicable
• or by a combination of the above.

A result mapping (implicit or explicit by a with clause) is well typed, if the value
at the left hand conforms to the right hand side according to the rules given above,
except that translation polymorphism here applies lowering instead of lifting.
These rules define translation polymorphism as introduced in §2.3.
Additionally, in a replace callin binding compatibility of parameters and return
types must hold in both directions. Thus, from the above list of conversions a
replace binding cannot apply subtype polymorphism nor primitive type conversion.
If more flexibility is desired, type parameters can be used as defined in §4.10.

(e) Role arrays

For arrays of roles as parameters §2.3.(d) applies accordingly. For arrays as return
value §2.2.(e) applies.

(f) Base calls

For base calls these rules are reversed again, i.e., a base call behaves like a callout
binding.

Overriding access restrictions §4.6.

Callin bindings may also mention inaccessible methods (cf. decapsulation §3.4.). Due to the
reverse call direction this is relevant only for base calls within callin methods. Base calls
have unrestricted access to protected base methods. Accessing a base method with private or
default visibility is also allowed, but signaled by a compiler warning.

Comment:
A base call to an inaccessible base method is considered harmless, since this is the originally
intended method execution.

(a) Private methods from super classes

(Cf. §3.4.(d)) If a callin binding shall bind to a private base method, that method
must be defined in the exact base class to which the current role class is bound using
playedBy.
If a private base feature must indeed be callin-bound, a role class must be defined
that is played by the exact base class defining the private feature. Another role
bound to a sub-base-class can then be defined as a sub class of the first role. It will
inherit the callin binding and through this it can access the desired feature.
§ 4.7. Callin binding with static methods

The normal case of callin bindings refers to non-static methods on both sides (base and role). Furthermore, in Java inner classes can not define static methods. Both restrictions are relaxed by the following rules:

(a) Static role methods
   A role class may define static methods (see also § 1.2.1.(f)).

(b) Binding static to static
   A callin binding may bind a static role method to one or more static base methods. It is, however, an error to bind a static base method to a non-static role method, because such binding would require to lift a base object that is not provided.

(c) before/after
   In addition to the above, before and after callin bindings may also bind a static role method to non-static base methods.

(d) replace
   In contrast to § 4.7.(c) above, a replace callin binding cannot bind a static role method to a non-static base method.

The following table summarizes the combinations defined above:

<table>
<thead>
<tr>
<th>role method</th>
<th>base method</th>
<th>static</th>
<th>non-static</th>
</tr>
</thead>
<tbody>
<tr>
<td>static</td>
<td>OK</td>
<td>before/after: OK</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>replace: illegal</td>
<td></td>
</tr>
<tr>
<td>non-static</td>
<td>illegal</td>
<td>OK</td>
<td></td>
</tr>
</tbody>
</table>

(e) No overriding

Since static methods are not dynamically bound, overriding does not apply in the normal semantics. Regarding callin bindings this has the following consequences (assuming a role RMid played by BMid plus its super-class BSuper and its sub-class BSub.

1. If a static base method BMid.m is bound by a callin binding this has no effect on any method m in BSub.

2. If a callin binding mentions a method m which is not present in BMid but resolves to a static method in BSuper the binding only affects invocations as BMid.m() but not BSuper.m(). If the latter call should be affected, too, the callin binding must appear in a role class bound to BSuper, not BMid.

3. In order to bind two static base methods with equal signatures, one being defined in a sub-class of the other one, two roles have to be defined where one role refines the playedBy clause of the other role (say: public class RSub extends RMid playedBy BSub). Now each role may bind to the static base method accessible in its direct base-class.
Callin precedence

Syntax § A.8

If multiple callins from the same team refer to the same base method and also have the same callin modifier (before, after or replace), the order in which the callin bindings shall be triggered has to be declared using a precedence declaration.

(a) Precedence declaration

A precedence declaration consists of the keyword precedence followed by a list of names referring to callin bindings (see § 4.1.(e) for named callin bindings).

```
precedence callinBinding1, callinBinding2;
```

A precedence declaration is only legal within a role or team class. The order of elements in a precedence declaration determines their priority during dispatch, similar to priorities based on activation of several team instances (§ 5.1). This means that before and replace binding with highest priority trigger first, whereas after bindings with highest priority trigger last. For binding precedences (as opposed to class based precedence, see § 4.8.(c) below) which refer to after bindings, the precedence declaration must also use the after keyword to remind the programmer that the execution order is inverse to the textual order.

```
precedence after importantExecuteLast, lessImportantExecuteEarlier;
```

(b) Qualified and unqualified names

Within a role class a callin binding may be referenced by its unqualified name. A precedence declaration in a team class must qualify the callin name with the name of the declaring role class. A team with nested teams may concat role class names. Elements of a qualified callin name are separated by ".".

The callin binding must be found in the role specified by the qualifying prefix or in the enclosing role for unqualified names, or any super class of this role (including implicit super classes § 1.3.1).

(c) Class based precedence

At the team level a precedence declaration may contain role class names without explicitly mentioning callin bindings in order to refer to all callin bindings of the role.

(d) Multiple precedence statements

All precedence statements are collected at the outer-most team. At that level all precedence declarations involving the same base method are merged using the C3 algorithm [3] (see p. 65). When merging precedences declarations more deeply nested declarations have higher priority than outer declarations. For several declarations at the same nesting level the lexical ordering determines the priority. At any point the C3 algorithm will ensure that the resulting order after merging is consistent with each individual precedence declaration. It is an error to declare incompatible precedence lists that cannot be merged by the C3 algorithm.
4. CALLIN BINDING

(e) Binding overriding

Precedence declarations may conflict with overriding of callin bindings (see §4.1.(e)):
For each pair of callin bindings of which one callin binding overrides the other one,
precedence declarations are not applicable, since dynamic binding will already select
exactly one callin binding.
It is an error to explicitly mention such a pair of overriding callin bindings in a prece-
dence declaration.
When a class-based precedence declaration implicitly refers to a callin binding that is
overridden by, or overrides any other callin binding within the same precedence decl-
aration, this does not affect the fact, that the most specific callin binding overrides
less specific ones.

Listing 18: Callin binding example

```java
public class LogLogin playedBy Database {
  callin void log (String what) {
    System.out.println("enter " + what);
    base.log(what.toLowerCase());
    System.out.println("leave " + what);
  }
  void log (String what) <- replace void login(String uid, String passwd)
    with { what <- uid }
  (new Database()).login("Admin", "Passwd");
}
```

Effects: Provided the callin bindings are active (cf. §5) then:

- the call in line 10 is intercepted by method log of role LogLogin.
- the call target of log is a role of type LogLogin which is created by lifting the original
call target (of type Database) to LogLogin.
- only parameter uid is passed to log (bound to formal parameter what).
- within method log the base call (line 4) invokes the original method passing a modified
  uid (converted to lower case, cf. line 4) and the unmodified password, which is hidden
  from the callin method due to the parameter mapping in line 8.

§ 4.9. Callin inheritance

This section defines how callin bindings and callin methods relate to inheritance.

§ 4.9.1. Base side inheritance

Generally, a callin binding affects all sub-types of its bound base. Specifically, if a role type R
bound to a base type B defines a callin binding rm <- callin_modifier bm, the following
rules apply:

(a) Effect on sub-classes

The callin binding also effects instances of any type BSub that is a sub-type of B.
If BSub overrides the bound base method bm, the overridden version is generally
affected, too. However, if \( \text{bm} \) covariantly redefines the return type from its super version, the callin binding has to explicitly specify if the covariant sub-class version should be affected, too (see \( \S \) 4.9.3.(b)).

### (b) No effect on super-classes

The binding never affects an instance of any super-type of \( B \) even if the method \( \text{bm} \) is inherited from a super-class or overrides an inherited method. This ensures that dispatching to a role method due to a callin binding always provides a base instance that has at least the type declared in the role’s playedBy clause.

For corresponding definitions regarding static methods see \( \S \) 4.7.(e).

#### Role side inheritance \( \S \) 4.9.2.

Any sub-type of \( R \) inherits the given callin binding (for overriding of bindings see \( \S \) 4.8.(e)). If the sub-role overrides the role method \( \text{rm} \) this will be considered for dynamic dispatch when the callin binding is triggered.

#### Covariant return types \( \S \) 4.9.3.

Since version 5, Java supports the covariant redefinition of a method’s return type (see JLS 8.4.5\(^2\)). This is not supported for callin methods (\( \S \) 4.9.3.(a)). If base methods with covariant redefinition of the return type are to be bound by a callin binding the subsequent rules ensure that type safety is preserved. Two constraints have to be considered:

1. When a callin method issues a base-call or calls its tsuper version, this call must produce a value whose type is compatible to the enclosing method’s declared return type.

2. If a replace-bound role method returns a value that is not the result of a base-call, it must be ensured that the return value actually satisfies the declared signature of the bound base method.

#### (a) No covariant callin methods

A method declared with the callin modifier that overrides an inherited method must not redefine the return type with respect to the inherited method. This reflects that fact that an inherited callin binding should remain type-safe while binding to the new, overriding role method. Binding a covariant role method to the original base method would break constraint (1) above.

#### (b) Capturing covariant base methods

If a callin binding should indeed affect not only the specified base method but also overriding versions which covariantly redefine the return type, the binding must specify the base method’s return type with a "+" appended to the type name as in

\[
\text{void rm()} \leftarrow \text{before \ RT+ bm()};
\]

Without the "+" sign the binding would only capture base methods whose return type is exactly \( \text{RT} \); by appending "+" also sub-types of \( \text{RT} \) are accepted as the declared return type.

\(^2\)http://java.sun.com/docs/books/jls/third_edition/html/classes.html#8.4.5
(c) **Covariant replace binding**

When using the syntax of §4.9.3(b) to capture base methods with covariant return types in a callin binding with the replace modifier, the role method must be specified using a free type parameter as follows:

\[
\langle E \text{ extends } RT \rangle E \text{ rm}() \leftarrow \text{replace } RT+ bm();
\]

The role method rm referenced by this callin binding must use the same style of return type using a type parameter. The only possible non-null value of type E to be returned from such method is the value provided by a base-call or a tsuper-call. This rule enforces the constraint (2) above.

Note that this rule is further generalized in §4.10.

**Listing 19: Binding a parametric role method**

```java
public class SuperBase {
    void check() { System.out.print("OK"); }
}

public class SubBase extends SuperBase {
    SubBase foo() { return this; }
    void print() { System.out.print("SubBase"); }
    String test() {
        this.foo().print(); // print() requires a SubBase
    }
}

public team class MyTeam {
    protected class R playedBy SuperBase {
        callin \langle E \text{ extends } SuperBase \rangle E ci() {
            E result = base.ci();
            result.check(); // check() is available on E via type bound SuperBase
            return result;
        }
        \langle E \text{ extends } SuperBase \rangle E ci() <- replace SuperBase+ foo();
    }
}
```

**Explanation:**

- Method `SubBase.foo` in line 7 redefines the return type from `SuperBase` (inherited version) to `SubBase`, thus clients like the method call in line 10 must be safe to assume that the return value will always conform to `SubBase`.
- The callin binding in line 21 explicitly captures both versions of `foo` by specifying `SuperBase+` as the expected return type. Thus, if an instance of `MyTeam` is active at the method call in line 10, this call to `foo` will indeed be intercepted even though this call is statically known to return a value of type `SubBase`.
- The callin method in lines 16-20 has a return type which is not known statically, but the return type is represented by the type variable `E`. Since the base call is known to have the exact same signature as its enclosing method, the value provided by the base call is of the same type `E` and thus can be safely returned from `ci`. Note, that no other non-null value is known to have the type `E`.
- By specifying `SuperBase` as an upper bound for the type `E` the callin method `ci` may invoke any method declared in type `SuperBase` on any value of type `E`. For an example see the call to `check` in line 18.
As an aside note that the above example uses type `SuperBase` in an undisciplined way: within role `R` this type is bound using `playedBy` and the same type is also used directly (as the upper bound for `E`). This is considered bad style and it is prohibited if `SuperBase` is imported using an base import (§ 2.1.2.(d)). Here this rule is neglected just for the purpose of keeping the example small.

### Generic callin bindings

As mentioned in § 4.5.(d) replace bindings do not support subtype polymorphism in either direction. On the other hand, binding several base methods to the same callin method may require some more flexibility if these base methods have different signatures. This is where type parameter come to the rescue to allow for generic callin methods and their binding to base methods with different signatures.

Note that this rule is a generalization of rule § 4.9.3.(c).

Additionally, any callin binding (before,replace,after) may declare one or more type parameters for propagating type parameters of the bound base method(s) (§ 4.10.(e)).

(a) **Fresh type parameter**

If a callin method declares a type parameter `<T>` for capturing a covariant return type this type `T` can be used for specifying the type of exactly one parameter or the return type. If a type parameter is used in more than one position of a callin method it is not considered a fresh type parameter and can thus not be bound to a covariant return type (see § 4.10.(d)).

(b) **Type bounds**

The type parameter of a callin binding may be bounded by an upper bound as in `<T extends C>`. In this case `T` can only be instantiated by types conforming to the upper bound `C`.

(c) **Generic replace binding**

A generic callin method according to the above rules is bound using a replace binding that declares the same number of type parameters, where type parameters of the binding and its callin method are identified. If the callin method declares bounds for its type parameters so should the replace binding.

(d) **Binding to a type parameter**

A fresh type parameter can be used to capture arbitrary types in the base methods to be bound. The type parameter may be instantiated differently for each bound base method. By such type parameter instantiation the types in role and base signatures are actually identical, thus satisfying the requirement of two-way substitutability.

Within the body of a generic callin method no further rules have to be followed, because the fresh type variable actually guarantees, that the role method cannot replace the original value (initial argument or base-call result) with a different object, because no type exists that is guaranteed to conform to the type parameters. Yet, the type bound allows the role method to invoke methods of the provided object.
Listing 20: Generic replace binding

```java
public team class MyTeam {
    protected class R playedBy Figures {
        callin <E extends Shape, F extends Shape> E ci(F arg) {
            E result = base.ci(arg);
            result = arg; // illegal, types E and F are incommensurable
            arg = result; // illegal, types E and F are incommensurable
            int size = arg.getSize(); // getSize() is available on F via type bound Shape
            result.resize(size); // resize() is available on E via type bound Shape
            return result; // only two legal values exist: result and null
        }
        <E extends Shape, F extends Shape> E ci(F arg) <- replace Rectangle getBoundingBox(Shape original) .
            Rectangle stretch(Square original);
    }
}
```

**Explanation:**

These declaration generate two version of the callin method ci:

1. Rectangle ci (Shape arg)
2. Rectangle ci (Square arg)

Within the callin method the following observations hold:

- Line 5 is illegal for the first signature as Shape is not conform to Rectangle
- Line 6 is illegal for the second signature as Rectangle is not conform to Square
- Everything else is type-safe.

**(e) Propagating type parameters**

If a callin binding binds to a generic base method, any type parameter(s) of the base method must be propagated into the role method by declaring the callin binding with type parameters, too. By matching a type parameter of a base method with a type variable of the callin binding, this genericity is propagated through the callin binding.

```java
class MyBase {
    <T> T getIt(T it) { return it; }
}
team class MyTeam {
    protected class MyRole playedBy MyBase {
        callin <U> U rm(U a) { return base.rm(a); }
        <U> U rm(U a) <- replace U getIt(U it);
    }
}
```

**Explanation:**
The callin binding declares a type parameter \(<U>\) which is used to match all occurrences of \(T\) in the signature of \(getIt\). Thus the implementation of \(rm\) uses the type \(U\) in exactly the same generic way as \(getIt\) uses \(T\).

Open issues:

The query language for specifying sets of base methods (§ 4.1.(d)) has not been finalized yet. In this version of the OTJLD § 8 acts as a placeholder for the section that will define a join point query language in the future.

References:

Team Activation

The concept of Activation

Binding activation  All callin bindings of a team only have effect if the team is active. Activation may be caused by explicit statements and also happens implicitly at certain points during program execution.

Guard predicates  Callin bindings can further be controlled using guard predicates, which can be attached to roles and teams at different levels. If a guard predicate evaluates to false, all affected callin bindings are disabled.

Effect of team activation

Activating a team instance has the effect of enabling all its callin bindings. All effects defined in § 4 apply only if a corresponding team instance is active.

The order of team activation controls the order of callin executions. If more than one team intercepts calls to the same base method, the most recently activated team has highest priority in that its before or replace callins are executed first while its after callins are executed last.

Global vs. thread local team activation

While thread local activation only enables the callin bindings of a team instance for a certain thread, global activation activates the team instance for all threads of the application.

Effect on garbage collection

Any active team is referenced by internal infrastructure. Thus, a team cannot be reclaimed by the garbage collector while it is active.

Explicit team activation

(a) Activation block

A team can be activated thread local by the block construct

```otjl
within (myTeam) { stmts }
```

If stmts has only one statement this can be abbreviated to

```otjl
within (myTeam) stmt
```

In these statements, myTeam must denote a team instance. For the time of executing this block, this team instance is activated for the current thread, which has entered the within block. The within block statement guarantees that it leaves the team in exactly the same activation state as it was in when entering this block. This includes the cases of exceptions, meaning that deactivation will also occur if the execution of the block terminates abnormally.
5. TEAM ACTIVATION

§ 5.2. Imperative activation

Each team class implicitly provides methods from the predefined interface org.objectteams.ITeam (super interface of all team classes) to control team activation disregarding the block structure of the program. The methods activate() and deactivate() are used to activate and deactivate a team instance for the current thread. If a team should be de-/activated for another thread this can be done by the methods activate(Thread aThread) and deactivate(Thread atThread). In order to achieve global activation for all threads the predefined constant org.objectteams.Team.ALL_THREADS is passed to these methods (e.g. activate(Team.ALL_THREADS)).

Note, that this methods make no guarantees with respect to exceptions.

§ 5.3. Implicit team activation

Implicit team activation is intended to ensure that whenever the control flow is passed to a team or one of its roles, the team is implicitly activated for the current thread. Implicit activation can be configured at different levels (see § 5.3.(d)).

When implicit activation is enabled a programmer may assume, that whenever a role forwards calls to its base object via callout, the callin bindings of the same role will be active at that time. Exceptions to this rule have to be programmed explicitly.

(a) Team level methods

While executing a team level method, the target team is always active. Activation is reset to the previous state when leaving the team method, unless the team has been explicitly activated during execution of the team method by a call to activate(). Explicit activation is stronger than implicit activation and thus persists after the team level method terminates. I.e., leaving a team level method will never reset an explicit activation.

(b) Methods of externalized roles

Invoking a method on an externalized role (see § 1.2.2) also has the effect of temporary activation of the team containing the role for the current thread. Regarding deactivation the rule of § 5.3.(a) above applies accordingly.

(c) Nested teams

Implicit activation has additional consequences for nested teams (see § 1.5):
Implicit activation of a team causes the activation of its outer teams.
Implicit deactivation of a team causes the deactivation of its inner teams.

(d) Configuring implicit activation

Implicit activation is disabled by default and can be enabled by adding the annotation `@org.objectteams.ImplicitTeamActivation`, which can be applied to a type or a method. When applied to a method it is ensured that invoking this method will trigger implicit activation. When the annotation is applied to a type this has the same effect as applying it to all externally visible methods of the type. Member types are not affected and have to be annotated separately.

The runtime environment can be configured globally by defining the system property `ot.implicit.team.activation` to one of these values:

- NEVER: Implicit activation is completely disabled.
- ANNOTATED: This is the default: implicit activation applies only where declared by `@ImplicitTeamActivation`.
- ALWAYS: Implicit activation applies to all externally visible methods (this was the default in OTJLD versions ≤1.2).

Note that among the different mechanisms for activation, `within` is strongest, followed by `(de)activate()`, weakest is implicit activation. In this sense, explicit imperative (de)activation may override the block structure of implicit activation (by explicit activation within a team level method), but not that of a `within` block (by deactivation from a `within` block).

Guard predicates

The effect of callins can further be controlled using so called guard predicates. Guards appear at four different levels:

- callin method binding
- role method
- role class
- team class

 Guards can be specified as regular guards or base guards, which affects the exact point in the control flow, where the guard will be evaluated.

(a) General syntax for guards

A guard is declared using the keyword `when` followed by a boolean expression in parentheses:

```
when (predicateExpression)
```
Depending on the kind of guard different objects are in scope using special identifiers like this, base.
Any predicate expression that evaluates to true enables the callin binding(s) to which it applies. Evaluation to false disables the callin binding(s).

(b) No side effects
A guard predicate should have no side effects. A compiler should optionally check this condition, but inter-procedural analysis actually depends on the availability of appropriate means to mark any method as side-effect free.

(c) Exceptions
A guard predicate should not throw any exceptions. Yet, any exception thrown within a guard predicate cause the guard to evaluate to false rather than propagating the exception, meaning that the evaluation of a guard predicate will never interrupt the current base behaviour.
A compiler should flag any checked exception that is thrown within a guard. Such diagnosis should by default be treated as an error, with the option of configuring its severity to warning or ignore.

§ 5.4.1. Regular guards
This group of guards evaluates within the context of a given role. These guards are evaluated after a callin target is lifted and before a callin bound role method is invoked.

(a) Method binding guards
A guard may be attached to a callin method binding as in:

```java
void roleMethod(int ir) <- after void baseMethod(int ib)
when (ir > MyTeam.this.threshold);
```

Such a guard only affects the callin binding to which it is attached, i.e., this specific callin binding is only effective, if the predicate evaluates to true.
The following values are within the scope of the predicate expression, and thus can be used to express the condition:
- The role instance denoted by this.
  Features of the role instance can also be accessed relative to this with or without explicit qualifying this.
- The team instance denoted by a qualified this reference as in MyTeam.this.
- If the callin binding includes signatures (as in the example above): Parameters of the role method.
  If parameter mappings are involved, they will be evaluated before evaluating the guard.

(b) Method guards
A method guard is similar to a method binding guard, but it applies to all callin method bindings of this method.
A method guard is declared between the method signature and the method body:
(c) Role level guards

When a guard is specified at the role level, i.e., directly before the class body of a role class, it applies to all callin method bindings of the role class:

```java
protected class MyRole
    when (value > MyTeam.this.threshold)
    {
        int value;
        other class body declarations
    }
```

The following values are within the scope of the predicate expression:

- The role instance denoted by `this` (explicit or implicit, see above). Thus, in the example `value` will be interpreted as a field of the enclosing role.
- The team instance denoted by a qualified `this` reference as in `MyTeam.this`.

(d) Team level guards

A guard specified in the header of a team class may disable the callin bindings of all contained role classes. The syntax corresponds to the syntax of role level guards. The only value directly available within team level guard is the team instance (denoted by `this`) and its features.

Of course all guards can also access any visible static feature of a visible class.

Even if a guard has no direct effect, because, e.g., a role class has no callin bindings (maybe not even a role-base binding), predicates at such abstract levels are useful, because all predicates are inherited by all sub classes (explicit and implicit).

▶ Base guards

The intention behind base guards is to prevent lifting of a callin-target if a guard evaluates to `false` and thus refuses to invoke the callin bound role method. Using base guards it is easier to prevent any side-effects caused by a callin binding, because lifting could cause side-effects at two levels:

- Creating a role on-demand already is a side-effect (observable e.g. by the reflective function `hasRole` (§ 6.1.)
- Role creation triggers execution of a role constructor (see `custom lifting constructor` (§ 2.3.1(c))) which could produce arbitrary side-effects.

Both kinds of side-effects can be avoided using a base guard which prevents unnecessary lifting.

Any guard (5.4.1 (b)-(e)) can be turned into a base guard by adding the modifier `base` as in:
protected class MyRole playedBy MyBase
    base when (base.value > MyTeam.this.threshold)
    {
        class body declarations
    }

However, different scoping rules apply for the identifiers that can be used in a base guard:

(a) Base object reference
   In all base guard predicates the special identifier base can be used to denote the base object that is about to be lifted.

(b) Method binding guards
   A base method binding guard may access parameters as passed to the base method. Parameter mappings are not considered. Additionally, for after callin bindings, the identifier result may be used to refer to the result of the base method (if any).

   Note: In order to achieve the same effect of accessing the base method’s result, a regular binding guard (not a base guard) must use a suitable parameter mapping (see §4.4.(c)).

(c) Method guards
   In contrast to regular method guards, a base guard attached to a role method cannot access any method parameters. See the next item (d) for values that are actually in scope.

(d) Role level guards
   Role level base guards may use these values:
   • The base instance using the special identifier base.
   • The team instance using a qualified this references (MyTeam.this).

(e) Team level guards
   Team level base guards have the same scope as role level base guards (d). However, the type of the role instance is not known here, i.e., here base has the static type java.lang.Object.

(f) Unbound roles
   In contrast to regular guards, base guards cannot be attached to unbound role classes nor to their methods. Only team level base guards are independent of role binding.
Multiple guards

Due to the different ranges of applicability different guards may affect the same method binding. In that case all applicable guards are conjoined using a logical and. Any guard is interpreted as the conjunction of these predicates (if present):

- The direct predicate expression of the guard.
- The next outer guard along the chain method binding -> method -> role level -> team level
- The guard at the same level that is inherited from the implicit super role.
- The guard at the same level that is inherited from the explicit super role.

Listing 21: Example code (Guard Predicates):

```java
public team class ATM {
    private Bank myBank;

    public class ForeignAccount playedBy Account {
        base when (ATM.this.myBank.equals(base.getBank())) {
            callin void debitWithFee(int amount) {
                base.debitWithFee(fee+amount);
            }
        }
        void debitWithFee(int i) <- replace void debit(int amount) {
            base when (amount < 1000);
        }
    }
}
```

Effects:

The team in this example causes that an additional fee has to be payed while debiting less than 1000 Euros from a "foreign" account.
The base guard in line 4 ensures that Account objects only get ForeignAccount roles, if they belong to a different bank than the surrounding ATM team. It accesses the bank of the base via the base identifier.

The method binding guard in line 10 restricts the callin to debitWithFee to calls where the base method argument amount is lower than 1000.

A call to Account.debit causes a replace callin to debitWithFee only if both predicates evaluate to true.

§ 5.5. Unanticipated team activation

If an application should be adapted unanticipatedly by one or more teams, this can be achieved without explicitly changing the program code of this application.

General activation via config file:
Instead of adding the team initialization and activation code to the main program, it is possible to add the respective teams via a config file. Every line of this text file contains the fully qualified name of a compiled team, which has to be available on the classpath. For the instantiation of these teams the default constructor is used, which means adding a team to an application this way requires the team to provide a default constructor. The activation order (see § 5.1) for these teams corresponds to the order in which they are listed in the config file. Lines starting with a ‘#’ denote comment lines.

Listing 22: Example config file:

```java
# Config file for an ObjectTeams application:
mypackage1.MyTeam1
# ...
mypackageM.MyTeamN
```

To get this config file recognized by the application the VM argument `-Dot.teamconfig= <config_file_name>` has to be used when starting the application.

Note:
In the ObjectTeams Development Tooling (OTDT) teams are activated unanticipatedly via a special tab in the “Run-Configuration” (see OTDT features\(^\text{21}\)), instead.

Activation adjustment example:
Teams added via the config file mechanism are activated by default. Because no reference to them is stored anywhere, it is not possible to deactivate them later. If deactivation of unanticipated added teams is required, this can be achieved by adding a manager team via config file and encapsulate the actual functionality in another team managed by the manager team. This way a functional team can be activated and deactivated as needed.

\(^{21}\)http://www.objectteams.org/distrib/features.html#execution
5.5. Unanticipated team activation

Listing 23: Example code (Activation Adjustment):

```java
public team class MyManagerTeam {
    private FunctionalTeam myFunctionalTeam = new FunctionalTeam();
    protected class MyRole playedBy MyApplication {
        void startAdaption() { myFunctionalTeam.activate(); }
        startAdaption <= before startMethod;
        void stopAdaption() { myFunctionalTeam.deactivate(); }
        stopAdaption <= after stopMethod;
    }
}
```

Effects:

- `startMethod` and `stopMethod` are methods which demand the activation and deactivation respectively.
- If the activation/deactivation depends on other conditions these can be checked in addition.
The role of predefined types and methods

Some features of OT/J are supported without introducing new syntax but by predefined types and methods.

Reflection

Object Teams supports reflection with respect to teams, roles, and role-base relationships.

(a) Interface to the role registry

Each team instance internally has a registry of known role objects indexed by their base object. Programmers may make use of this registry using the following reflective methods defined in org.objectteams.ITeam:

- `boolean hasRole (Object aBase)`; This method checks whether a role for the passed base object already exists in the target team.
- `boolean hasRole (Object aBase, Class roleType)`; This method checks whether an instance of type `roleType` as a role for the passed base object `aBase` already exists in the target team. The role may also be of any subtype of the specified role type. If `roleType` is not a member type of the current team an `IllegalArgumentException` is thrown.
- `Object getRole (Object aBase)`; If the passed base object `aBase` already has a role in the target team, this role is returned. Otherwise `null` is returned.
- `<T> T getRole (Object aBase, Class<T> roleType)`; If the passed base object `aBase` already has a role in the target team that is assignable to the type represented by `roleType`, this role is returned. Otherwise `null` is returned. If `roleType` is not a member type of the current team an `IllegalArgumentException` is thrown.
- `Object[] getAllRoles ()`; Retrieves all existing (registered) bound roles in the target team. This method uses internal structures of weak references. For that reason it may return role instances which were about to be reclaimed by the garbage collector. If performance permits, it is thus advisable to always call `System.gc()` prior to calling `getAllRoles()` in order to achieve deterministic results (see also § 2.1.(f)).
- `<T> T[] getAllRoles (Class<T> roleType)`; Retrieves all existing (registered) bound roles in the target team that are assignable to the type represented by `roleType`. If `roleType` is not a member type of the current team an `IllegalArgumentException` is thrown. See the note about garbage collection above.
- `void unregisterRole (Object aRole)`; This method unregisters the passed role object from the target team. Thus the corresponding base looses this role. After calling this method the role should no longer be used.
void unregisterRole ( Object aRole, Class roleType ) ; This method un-
registers the passed role object from the target team. Thus the corresponding
base loses this role. After calling this method the role should no longer be used.
The only difference to the previous method is improved speed because no search
for the corresponding registry has to be performed.
If roleType is not a member type of the current team an IllegalArgumentException
is thrown.

It is desirable and possible to use these methods within guards (see § 5.4). These
methods allow to write the specification of guards in a more concise and more
expressive way. Determined by the signature, the first four methods can only be
used in a base-level guard (§ 5.4.2) because they require a reference to a base
object.

Listing 24: Example code (Guards and Reflection):

```java
public team class SpecialConditions {
    public void participate(Account as BonusAccount ba) {}
    public class BonusAccount playedBy Account
        base when ( SpecialConditions.this.hasRole(base, BonusAccount.class) )
        { callin void creditBonus(int amount) {
            base.creditBonus(amount + bonus);
        }
        void creditBonus(int amount) <- replace void credit(int i)
            base when (i > 1000);
    }
}
```

Effects:
This teams provides a bonus system for registered Accounts. Every time an amount
of more than 1000 is deposited to a registered account, additional 1% of the amount
is credited.

- The team level method participate in line 2 uses declared lifting (see § 2.3.2) to allow the passed Account object to participate the bonus system provided by the SpecialConditions team.
- The base guard in line 4 uses the reflective method hasRole to check whether the base object already has a role of type BonusAccount in the surrounding team. The expression BonusAccount.class returns the java.lang.Class object representing the role BonusAccount (see JLS §15.8.22). This guard ensures that only accounts explicitly registered via participate are ever decorated with a role of type BonusAccount.
- The method binding guard in line 10 restricts the callin to creditBonus to calls where the base method argument amount is greater than 1000.

(b) Behavioral reflection
The following reflective methods defined in org.objectteams.ITeam can be used to inspect the dynamic behavior of a team:

```java
boolean isExecutingCallin () ; This method is used to inspect whether a
control flow has already been intercepted by at least one callin binding of the
current team. It can be used to avoid undesirable re-entrance to a team.
```
boolean isActive () ; This method checks whether the team instance is active for the current thread.

boolean isActive ( Thread aThread ) ; This method checks whether the team instance is active for the thread aThread.

(c) Class literals for roles

The Java syntax for so-called class literals, MyClass.class (see JLS §15.8.2\(^2\)) can be used for role types with slightly changed semantics: Role types are virtual types (§1.3.1.) that are bound dynamically §1.3.1.(e). This applies to role class literals, too. From this follows the constraint that a role class literal can only be used within the non-static context of a team, i.e., for evaluating a role class literal an enclosing team instance must be in scope.

Unlike regular type checking for role types, the class literal itself does not have a dependent type. Thus type checking of calls to methods like hasRole(Object, Class) cannot detect, whether the Class instance has actually been obtained from the correct team instance. Any attempt to pass a class that is not known as a bound role within the given team results in an IllegalArgumentException at run-time.

**Other API Elements** § 6.2.

(a) Interfaces for role encapsulation

A set of pre-defined types exist that do not extend java.lang.Object and have no features except the operators == and !=.

Note: The JLS defines that each interface declares all methods defined in java.lang.Object (JLS §9.2\(^2\)) and also each object referenced by an interface type can be widened to java.lang.Object. Compilers commonly implement this by declaring java.lang.Object the super-type of all interfaces. Such implementation has no visible difference with respect to the more complex definition in the JLS.

These predefined types are

org.objectteams.IConfined regular interface
org.objectteams.ITeam.IConfined role interface
org.objectteams.Team.Confined role class

These types provide no new functionality but inheriting from these types influences the semantics with respect to encapsulation. The purpose and usage of these types is described in § 7.

(b) Interface for explicit lowering

The following role interface exists for the purpose of allowing explicit lowering:

org.objectteams.ITeam.ILowerable role interface

This interface was introduced in detail in § 2.2.(d)\(^2\)

(c) Team activation methods

Every team can be activated and deactivated by predefined methods of the interface org.objectteams.ITeam.

activate() and activate(Thread th) Methods for activation of a team

deactivate() and deactivate(Thread th) Methods for deactivation of a team

The usage of these Methods is described in § 5.2.(b).

(d) Exceptions

The following Exceptions can be thrown during the execution of an ObjectTeam/Java program:

ResultNotProvidedException Thrown if a replace call in without a base call does not provide the necessary (primitive type) base result (see § 4.3.(e)).

LiftingFailedException Thrown if an actual ambiguity occurs during lifting (see § 2.3.4.(c)) or if lifting would need to instantiate an abstract role class (see § 2.5.(b)). This is a checked exception. See § 2.3.5. for more information.

WrongRoleException Thrown during lifting if the base object has, with respect to the same team instance, previously been lifted to a role type that is not conform to the currently requested type (see § 2.3.4.(d) and § 2.4.3.).

DuplicateRoleException Thrown during explicit role creation, if a new role is created for a base object, which already has a role of the required type in the given team (see § 2.4.1.(c)).

RoleCastException Thrown during cast of an externalized role, if the casted expression is anchored to a different team instance than the cast type (see § 1.2.4.(b)).

LiftingVetoException This exception is used internally to abort the process of lifting when a relevant guard predicate (§ 5.4.) evaluated to false. Such exceptions thrown from generated code will never appear in client code, so there is usually no need to catch a LiftingVetoException. However, in some situations it is useful to explicitly throw a LiftingVetoException from a lifting constructor (§ 2.3.1.(b)) of a role. This style allows to abort lifting even after the lifting constructor has started to work and also for method parameters requiring lifting. If lifting was triggered due to a call in method binding, this binding will simply not trigger if a LiftingVetoException is thrown while preparing the call to the role method.

(e) Role migration

The following interfaces can be used to enable role migration:

IBaseMigratable This interface declares a method

\[ \text{<茬> void migrateToBase(B otherBase)} \]

and instructs the compiler to generate an implementation of this method for any bound role declaring IBaseMigratable as its super-interface.

The effect of calling migrateToBase on a role instance is to re-bind this role to a new base instance. The base instance must be compatible to the role's
base class (in order to avoid problems during lifting the compiler may require the base to be of the exact type of the role's base class). Passing null to this method causes a NullPointerException to be thrown.

**ITeamMigratable** This interface declares a method

```java
<R> R<@otherTeam> migrateToTeam(final ITeam otherTeam)
```

and instructs the compiler to generate an implementation of this method for any role declaring `ITeamMigratable` as its super-interface. The effect of calling `migrateToTeam` on a role instance is to re-bind this role to become a contained part of a new team instance. The team instance must be of the exact type of the role's enclosing team. Passing null to this method causes a NullPointerException to be thrown.

**Caveat:**

This method intentionally breaks the rules of family polymorphism: any reference `R<@previousTeam> r` which was established before migration will incorrectly imply that the role's enclosing team still is `previousTeam`, which is no longer true after migration. While this does not effect any method lookup (which is still safe), further assumptions based on a role's dependent type are invalidated by team migration. The same holds for references from the migrating role to any sibling role instances. If the rules of family polymorphism should be maintained one should just refrain from declaring `ITeamMigratable` as a role's super-interface.

For both methods the signature declared in the interface is over-generalized, yet the compiler performs the necessary checks to ensure that role, base and team instances are indeed compatible and additionally the return type of `migrateToTeam` is checked as a self-type, i.e., it reflects the exact type of the call target.

---

### Annotations

#### (a) Controlling implicit team activation

Implicit team activation is disabled by default and can be enabled by adding the annotation @org.objectteams.ImplicitTeamActivation. See §5.3.(d) for details.

#### (b) Controlling lifting

If lifting as defined in §2.3 and specifically §2.3.1 causes performance problems, the semantics of lifting can be modified per role class using the annotation @org.objectteams.Instantiation. See §2.3.1.(d) for details.
§ 7. Role Encapsulation

Concepts of encapsulation

Protected roles A role with visibility protected cannot be externalized, which means its type cannot be used outside the declaring team (§ 1.2.3).

Confined roles Confined roles are encapsulated even stricter than protected roles: the compiler will ensure that by no means any object outside the enclosing team will ever have a reference to a confined role.

Opaque roles Opaque roles build on the guarantees of confined roles but allow to be shared in a limited way such that no information is exposed.

§ 7.1. Opaque roles

The purpose of the two IConfined interfaces (see § 6.2.(a)) is to define opaque roles: Any role implementing IConfined can be externalized using this type, such that external clients cannot access any features of the role. The type IConfined exposes no features and references of this type cannot be widened to any type not even to java.lang.Object.

If the actual role type is furthermore invisible outside the team (by not declaring it public), it is perfectly safe to externalize such roles using type IConfined (which is a public interface) and pass them back to the owning team. The encapsulation of the team is in no way breached by externalizing opaque roles, which can only be used as a handle into internal state of the team.

The difference between the two mentioned interfaces is that ITeam.IConfined requires to use this type or any subtype as externalized role. Such a reference contains the information of the enclosing team. Even stricter control can be imposed using the regular interface IConfined. Here not even team membership is visible to clients using a reference of this type.

§ 7.2. Confined roles

Subclassing Team.Confined with a protected class yields a role class to which no object outside the team will ever have a reference. The point here is that instances of a role class with a regular super class can be widened to this super class. Widening can occur either in an assignment or when invoking a method which the role inherits from the regular super class, where the this reference is widened. In both cases the widened reference is no longer protected by the team and can leak out. This would break encapsulation of a role object that should only be accessible within the enclosing team.

Subclasses of Team.Confined are not compatible to any class outside their enclosing team (including java.lang.Object) and do not inherit any methods that have the danger of leaking this.

(a) Inhibition of overriding

The types ITeam.IConfined and Team.Confined cannot be overridden (cf. § 1.3.1.(c)).

(b) Arrays of Confined

For any confined type C, i.e., a type which is not compatible to Object, an array of C is not compatible to an array of Object nor to Object itself. This rule ensures
that confinement cannot be bypassed by a sequence of compatible assignments and casts.

**Upcoming:**

*Only by widening to a non-role super-type, a role instance can be accessed from outside the team. In the future this can be inhibited by restricted inheritance.*

Listing 25: Example code (Role Encapsulation):

```java
public team class Company {
    private HashMap<String, Employee> employees;

    protected class Employee implements IConfined {
        void pay(int amount) { ... }
    }

    public IConfined getEmployee(String ID) {
        return employees.get(ID);  // implicit widening to IConfined
    }

    public void payBonus(IConfined emp, int amount) {
        ((Employee)emp).pay(amount);  // explicit narrowing
    }
}

public class Main {
    public static void main(String[] args) {
        final Company comp = new Company();
        IConfined<@comp> emp = comp.getEmployee("emp1");
        // System.out.println(emp); <- forbidden!
        comp.payBonus(emp, 100);
    }
}
```

**Effects:**

- The protected role Employee implements the above described interface IConfined and therefore becomes opaque (line 4).
- Methods sharing such an opaque role with the outside of the enclosing team have to use the type IConfined (line 8, line 11).
- It is possible to obtain an instance of such a role by using the type IConfined (line 18).
- Trying to access any feature of this instance, for example toString(), will cause a compilation error (line 19).
- Passing the opaque role reference back into the team works well (line 20).
- Inside the team some conversions between the types IConfined and the intrinsic role type Employee may be necessary (line 9 and 12).
Join Point Queries

Defining sets of join points for interception

join point  In OT/J a join point is considered to be an element of the program. A meta model exists which defines the kinds of join points that can be identified.

join point interception  The purpose of identifying join points is to intercept program execution at these points by means of calling bindings.

join point query  Sets of join points are defined using functional queries of the program’s reflective representation.

pointcuts  Dynamic "pointcuts" comparable to AspectJ’s cflow or even the Trace-Matches approach are not subject to the join point language of OT/J. These features will be added at a different level of abstraction. Note, that guard predicates (§ 5.4.) subsume the dynamic capabilities of the pointcuts if, target, this, args.

Join point queries

This section will describe the query language used to define sets of join points. As of version 1.4.0 of the OTDT this query language is not yet supported.

Query expressions

OT/J meta model
Value Dependent Classes

Generalizing externalized roles

Type Value Parameter In addition to regular generics, a class may declare parameters that represent an object value. Such a value parameter is called the type anchor for this class, the class’s type is said to be anchored to this parameter.

Value Dependent Classes A class that declares one or more value parameters depends on the runtime instance(s) denoted by its anchor(s).

Externalized Roles The concept of externalized roles (§ 1.2.2) is a special case of the concepts presented here.

Defining classes with value parameters

(a) Value parameter declaration

Within the angle brackets that mark the parameters of a generic class also value parameters can be declared. In contrast to a type parameter, a value parameter is denoted as a pair of two identifiers: a type and a free name, e.g.,

\[
\text{class MyClass<YourType aName> \{ \ldots \}}
\]

Note that value parameters are valid for classes only, not for interfaces.

(b) Value parameter application

Within the given class (MyClass) the parameter name (aName) can be used like a final field of the given type (YourType). In contrast to regular final fields the assignment to this name occurs even before the constructor is executed.

(c) Role types as dependent types

Any role type can be interpreted as a value dependent type, however, in the declaration of a role type the value parameter remains implicit: it is identical to the enclosing team instance.

Using classes with value parameters

When using a class which declares one or more value parameters (type anchors) a corresponding anchor value has to be provided.

Parameter substitution

Substitution of a type anchor of a class MyClass<YourType p> is denoted as MyClass@$v$. In this term $v$ must be a value which is conform to the declaration of the value parameter "YourType p", i.e., $v$ must have the static type YourType.
The value passed for substituting a type anchor must be a path of variables declared as \texttt{final}. Obviously, only the first element in such a path can be a local variable or a method argument, all other elements have to be fields. The reason for requiring final variables is in type checking as discussed next.

\textbf{Note:} 

Externalized roles as defined in \S 1.2.2.(b) are a special case of types with a value parameter, where the value is an instance of the enclosing team.

\textbf{(a) Instance constrained type parameters}

In addition to normal usage, a value parameter can be applied nested to a regular type parameter:

\begin{verbatim}
class MyClass<YourType aName, DependentParam<aName>> { . . .
\end{verbatim}

Here the type parameter \texttt{DependentParam} is constrained to be anchored to \texttt{aName}. If a value parameter is used as a constraint for a regular type parameter any substitution for the type parameter must also supply a value matching the value parameter. The class from above could be applied like this:

\begin{verbatim}
final YourType anchor = new YourType();
MyClass <@anchor, YourDependent<@anchor>>
\end{verbatim}

Within the declaring element (class or method) applications of the type variable representing the instance constrained type parameter must repeat the anchor verbatim, i.e., no substitutions are performed here.

\section{Type conformance}

Two value dependent types (anchored types) are considered conform only if the anchors of both types refer to the same object(s). The compiler must be able to statically analyze this anchor identity.

\textbf{(a) Substitutions for type anchors}

Only two substitutions are considered for determining anchor identity:

1. If a method signature uses \texttt{this} as the anchor of any of its types, type checking an application of this method performs the following substitutions:
   - A simple \texttt{this} expression is substituted by the actual call target of the method application.
   - A qualified \texttt{Outer.this} expression is substituted by the corresponding enclosing instance of the call target.

2. Assignments from a \texttt{final} identifier to another \texttt{final} identifier are transitively followed, i.e., if t1, t2 are final, after an assignment \texttt{t1=t2} the types \texttt{C<@t1>} and \texttt{C<@t2>} are considered identical. Otherwise \texttt{C<@t1>} and \texttt{C<@t2>} are incommensurable.
   - Attaching an actual parameter to a formal parameter in a method call is also considered as an assignment with respect to this rule.
(b) Conformance of raw types

After anchors have been proven identical, the raw types are checked for compatibility using the standard Java rules.

Restrictions and limitations

(a) No overriding

Types with value parameters that are declared outside a team cannot be overridden, as roles can be. Therefore, implicit inheritance does not apply for these types.

(b) Only first parameter

Currently only the first parameter of a class may be a value parameter. This restriction may be removed in the future.
OT/J Syntax

Notation

The following grammar rules extend the Java grammar given in the [Java Language Specification](http://java.sun.com/docs/books/jls/second_edition/html/syntax.doc.html). We adopt the conventions of printing non-terminal symbols in italic font (e.g., `ClassDeclaration`), and terminal symbols in roman font (e.g., `class`). Names printed in black refer to definitions from the original Java grammar. Object Teams additions are printed in [blue boldface](#). For those rules that simply add a new option to an existing rule, the original options are indicated by an ellipse (…).

Keywords

The keywords introduced by OT/J have different scopes, which means outside their given scope these keywords can be used for regular identifiers. Only these names are keywords unconditionally:

```plaintext
readonly, team, within
```

Scoped keywords

The following names are keywords in OT/J only if they appear within a team or role class, i.e., after the keyword `team` has been recognized:

```plaintext
as, base, callin, playedBy, precedence, tsuper, with, when
```

These names are keywords only in the context of a callin or callout binding respectively (§A.3):

```plaintext
after, before, replace, get, set
```

Inheriting scoped keywords

While regular Java classes may use the scoped keywords (§A.0.1) of OT/J freely, it is an error if a role class inherits a feature whose name is a scoped keyword.

Internal names

Compiler and runtime environment generate internal methods and fields which start with the prefix `_OT$`. It is illegal to use any of these methods and fields within client code.

Class definitions

Class definitions add two new keywords `team` and `playedBy`. Classes which use these keywords are called teams and bound roles, respectively. Any class that inherits from a bound role class (either by an `extends` clause or by implicit inheritance, cf. [§1.3.1(c)](http://java.sun.com/docs/books/jls/second_edition/html/expressions.doc.html)) is again a bound role class.
§ A.1.1 **ClassDeclaration:**  
[Modifiers] [team] class Identifier [extends Type] [implements TypeList]  
[playedBy Type] [Guard] ClassBody

Contextual constraints:

(a) A class which has a playedBy clause (a **bound role** class) may not be declared static and must be directly contained in a class that has the **team** modifier (a **team** class).

(b) A class which inherits from a team class must have the team modifier, too.

(c) A class which has a guard (see § 5.4) must be a team or a role.

§ A.2. **Modifiers**

The rule for method modifiers adds one keyword: **callin**:

§ A.2.1 **Modifier:**

...  

callin

Contextual constraints:

(a) The class of a method which has the callin modifier may not be declared static and must be directly contained in a team class.

(b) A method that has the callin modifier may not appear in an explicit method call (rule Apply in JLS).

§ A.3. **Method bindings**

The rule of items declarable in a class body is augmented by method bindings:

§ A.3.1 **ClassBodyDeclaration:**

...  

callinbinding

calloutbinding

§ A.3.2 **CalloutBinding:**

[Modifier] [TypeArguments] MethodSpec CalloutKind MethodSpec  
CalloutParameterMappings  
[Modifier] [TypeArguments] MethodSpec CalloutKind CalloutModifier  
FieldSpec

§ A.3.3 **Callin binding:**

[ Identifier : ] [TypeArguments] MethodSpec <- CallinModifier MethodSpec  
[Guard] CallinParameterMappings

§ A.3.4 **MethodSpec:**

Identifier  
ResultType MethodDeclarator

Note, that ResultType and MethodDeclarator are not explicit in the overall syntax of the Java language specification. For convenience we refer to the definition in section 8.4. Method Declarations of the Java language specification.

§ A.3.5 \textbf{MethodSpecs:}
\begin{itemize}
\item \texttt{MethodSpec}, \texttt{MethodSpecs}
\end{itemize}

§ A.3.6 \textbf{CalloutKind:}
\begin{itemize}
\item \texttt{->}, \texttt{=>}
\end{itemize}

§ A.3.7 \textbf{CallinModifier:}
\begin{itemize}
\item \texttt{before}
\item \texttt{after}
\item \texttt{replace}
\end{itemize}

§ A.3.8 \textbf{CalloutModifier:}
\begin{itemize}
\item \texttt{get}
\item \texttt{set}
\end{itemize}

§ A.3.9 \textbf{FieldSpec:}
\begin{itemize}
\item \texttt{[Type]} \texttt{Identifier}
\end{itemize}

Contextual constraints:
\begin{itemize}
\item (a) CalloutBindings and CallinBindings may occur only in bound role classes.
\item (b) A CalloutBinding or CallinBinding may not mix identifiers and full signatures (MethodDeclarationHead) for its method specifiers (MethodSpec).
\item Binding a full method signature to a field requires the FieldSpec to include the \texttt{Type}.
\item (c) The method specifier at the left hand side of a CallinBinding which has the \texttt{replace} modifier must refer to a method that has the callin modifier.
\item (d) The \texttt{Modifier} of a callout binding can only be one of the visibility modifiers public, protected or private. A short callout binding (i.e., without signatures) must not specify a visibility modifier.
\end{itemize}

Parameter mappings
\textbf{§ A.4.}

§ A.4.1 \textbf{CalloutParameterMappings:}
\begin{itemize}
\item \texttt{with} \{ \texttt{CalloutParameterMappingList \[\]} \}
\end{itemize}

§ A.4.2 \textbf{CallinParameterMappings:}
\begin{itemize}
\item \texttt{with} \{ \texttt{CallinParameterMappingList \[\]} \}
\end{itemize}

§ A.4.3 \textbf{CalloutParameterMappingList:}
\begin{itemize}
\item \texttt{CalloutParameterMapping \[\]}, \texttt{CalloutParameterMappingList}
\end{itemize}

§ A.4.4 \textbf{CallinParameterMappingList:}
\begin{itemize}
\item \texttt{CallinParameterMapping \[\]}, \texttt{CallinParameterMappingList}
\end{itemize}

§ A.4.5 \textbf{CalloutParameterMapping:}
\begin{itemize}
\item \texttt{Expression \texttt{-\textgreater} Identifier}
\item \texttt{result \texttt{<-} Expression}
\end{itemize}

§ A.4.6 \textbf{CallinParameterMapping:}
\begin{itemize}
\item \texttt{Identifier \texttt{<-} Expression}
\item \texttt{Expression \texttt{-\textgreater} result}
\end{itemize}

\textbf{Note:}
\begin{itemize}
\item By defining ",;" as an option for parameter mappings, the grammar enforces that method
A.5. Statements

\[\text{Statement:}\]
\[\text{Within}\]
\[\text{BaseCall}\]
\[\text{TSuperCall}\]

A.5.1

\[\text{Within:}\]
\[\text{within ( Expression ) Statement}\]

A.5.2

\[\text{BaseCall:}\]
\[\text{base . Identifier ( Arguments\textsubscript{opt} )}\]
\[\text{base ( Arguments\textsubscript{opt} )}\]

A.5.3

\[\text{TSuperCall:}\]
\[\text{tsuper . Identifier ( Arguments\textsubscript{opt} )}\]
\[\text{tsuper ( Arguments\textsubscript{opt} )}\]

Contextual constraints:

(a) The expression of a Within must evaluate to an instance of a team class.

(b) The first form of a BaseCall may occur only in the body of a method that has the callin modifier. The identifier must be the name of the enclosing method.

(c) The second form of a BaseCall may occur only in a constructor of a bound role class.

(d) The first form of a TSuperCall may occur only in a method of a role class.

(e) The second form of a TSuperCall may occur only in a constructor of a role class.

A.6. Types

\[\text{Type:}\]
\[\text{LiftingType}\]
\[\text{AnchoredType}\]

A.6.1

\[\text{LiftingType:}\]
\[\text{Type as Type}\]

A.6.2

\[\text{AnchoredType:}\]
\[\text{Path . Type}\]

A.6.3

\[\text{Path:}\]
\[\text{Identifier}\]
\[\text{Path . Identifier}\]

Contextual constraints:

(a) Location

A LiftingType may only occur in the parameter list of a method of a team class.
(b) Role in scope
The right hand side type in a LiftingType must be a class directly contained in the enclosing team class (the class may be acquired by implicit inheritance (§ 1.3.1(c))).

(c) Team path
Note, that the syntax of §A.6.3/4 is deprecated in favor of § A.9.
The path in an AnchoredType must refer to an instance of a team class. Each identifier in the path must be declared with the final modifier.

Guard predicates

§ A.7.1 Guard:
\[ \text{base} \text{ when ( Expression )} \]

§ A.7.2 MethodDeclaration:
\[ \ldots \\ \text{MethodHeader \{Guard\} MethodBody} \]

Other rules referring to Guard:
ClassDeclaration (§ A.1.) CallinBinding (§ A.3.)

Contextual constraints:
(a) The Expression in a guard must have type boolean.

Precedence declaration

§ A.8.1 PrecedenceDeclaration:
\[ \text{precedence \{after\} CallinNameList ;} \]

§ A.8.2 CallinNameList:
\[ \text{Name [, CallinNameList]} \]

Value dependent types

§ A.9.1 TypeParameter:
\[ \text{TypeVariable \{TypeBound\} ReferenceType Name} \]

See JLS 3 §4.4

§ A.9.2 ActualTypeArgument:
\[ \text{ReferenceType Wildcard \@Name} \]

See JLS 3 §4.5.1

Contextual constraints:
(a) ActualTypeParameter
An ActualTypeArgument of the form \@Name may only occur as a parameter of a simple name type reference.

\[ \text{http://java.sun.com/docs/books/jls/third_edition/html/typesValues.html#108850} \]
\[ \text{http://java.sun.com/docs/books/jls/third_edition/html/typesValues.html#107353} \]
§ A.10. Packages and imports

<table>
<thead>
<tr>
<th>§ A.10.1 PackageDeclaration:</th>
</tr>
</thead>
<tbody>
<tr>
<td>...</td>
</tr>
<tr>
<td>team QualifiedName ;</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>§ A.10.2 Import:</th>
</tr>
</thead>
<tbody>
<tr>
<td>...</td>
</tr>
<tr>
<td>import base QualifiedName ;</td>
</tr>
</tbody>
</table>
Changes between versions

Paragraphs changed between versions (a) Between OTJLD 1.0 and OTJLD 1.1

- § 3.2.(a): Parameter mappings
  Disallow parameter mappings in a role interface.
- § 4.5.(d): Replace bindings
  Disallow unsafe use of polymorphism and primitive type conversions.
- § 6.1.(a): Signatures of reflective methods
  Made two methods generic so that return values can be used without the need of casting.
- § 7.2.: Confined roles
  Improved explanation.

(b) Between OTJLD 1.1 and OTJLD 1.2

- § 1.2.1.(e): Visibility of role features
  Clarification has been added that a role can always access all the features that its enclosing team has access to.
- § 2.1.2.(e): Generic roles/bases
  Relaxed the rules about generic bound roles. This change also subsumes what previously was a specific restriction in § 4.1.(b)
- § 3.1.(i) and § 3.5.(f): Visibility of shorthand callout
  A role method defined by a shorthand callout binding can now specify a visibility modifier (see also § A.3.), otherwise it inherits the visibility modifier of its bound base method/field.
- § 3.1.(j) and § 3.5.(h): Visibility of inferred callout
  Role methods inferred as a callout binding are either public (inferred via interface) or private inferred from self call/field access.
- § 3.5.(h): No explicit use of inferred callout to field
  Clarification has been added that an accessor method generated for an inferred callout to field can not be explicitly invoked.
- § 4.1.(b): No callin in generic role
  A restriction has been made explicit that a generic role cannot define callin bindings.
- § 4.2.(d): Callin methods
  Slightly rephrased and extended the rule to make explicit that a callin method can indeed be intercepted using a second level callin binding.
- § 6.1.(a): Reflective methods getAllRoles
  More precision: answer only bound roles.
(c) Between OTJLD 1.2 and OTJLD 1.3

- **§ 1.2.4.(c): Syntax for role class literals**
  Previously, the syntax `R<@t>.class` was not supported. This restriction has been removed.

- **§ 1.3.: Teams extending non-team classes**
  Previously, `org.objectteams.Team` was the super class of all team classes. As a consequence a team could not extend a non-team class. This restriction has been removed by introducing a new super-type of all teams, the interface `org.objectteams.ITeam`. This change also affects some paragraphs in § 6, as members have been moved to the new interface.

- **§ 1.5.(e): Precedence among different implicit supers**
  Corrected an inconsistency in the rules for precedence among different supers: The primary rule has always been that implicit inheritance binds stronger than explicit inheritance, however, for precedence among different implicit supers a different rule was defined. This has been changed such that different implicit supers are prioritized by the precedence of their enclosing teams, such that a role from an implicit super team is closer related than a role from an explicit super team.

- **§ 2.1.2.(b): Relaxed the rule against base class circularity**
  Base class circularity as defined in § 2.1.2.(b) is no longer an error but as a configurable warning. However, in the presence of base class circularity, neither callouts (§ 3.1.(a)) nor base constructor calls (§ 2.4.2.) are allowed.

- **§ 2.3.4.: Changed handling of role binding ambiguities**
  A definite binding ambiguity is no longer a (suppressable) compiler error, but is signaled by the need to declare `org.objectteams.LiftingFailedException`. This way diagnostics could be moved from rather unspecific locations in the team towards those applications that could suffer at runtime from a lifting failure. While it is generally not recommended to ignore any `LiftingFailedException` catching this exception may still make sense in a few corner cases mentioned in § 2.3.4.(b).

- **§ 4.4.(c): Further restrict result mapping in after callin bindings**
  Clarify that after callin bindings cannot use the `->` token to map a result value.

- **§ 4.8.(a): Precedence declarations affecting after callin bindings.**
  While previously the effect of precedence declarations was underspecified it has been defined that the order of elements in a precedence declaration affects their priority similar to § 5.1. This implies that the execution order for after bindings is now reversed compared to the previous implementation. In order to visualize this in the program it is now mandatory to mark precedence declarations for after bindings with the keyword `after`.

- **§ 4.10, § 4.10.(a): Generic callin bindings**
  Minor changes to give room for new paragraph § 4.10.(e)

- **§ 5.4.1.(a): Scope of regular binding guard**
  Removed an erroneous sentence about the special identifier `result` in a regular method binding guard. Since parameter mappings are applied before evaluating
the guard, the result value can be accessed through a result mapping (§ 4.4.(c)).
Furthermore, the sentence actually confused base and role sides.

- **§ A.3.** Syntax: generic method bindings
  The location of possible type parameters in a method binding has been made explicit.

### Additions between versions

**§ B.2.**

**(a) Between OTJLD 1.0 and OTJLD 1.1**

- **§ 1.2.4.(c):** Role class literal
  Made existing feature explicit and introduce new qualified class literal for externalized roles.

- **§ 3.1.(j) and § 3.5.(h):** Inferred callout
  New feature.

- **§ 4.6.(a):** Callin-binding private methods from super classes
  Added a necessary restriction.

- **§ 4.9:** Callin inheritance
  Clarified issues that where under-specified or insufficiently explained, specifically:
  - Effect of callin bindings on inherited or overridden base methods (§ 4.9.1.).
  - Interplay of callin bindings and base methods with covariant return types (§ 4.9.3.).

- **§ 4.10:** Generic replace bindings
  Reconcile type safety of replace bindings as introduced in § 4.5.(d) with desirable flexibility by using type parameters.

- **§ 7.2.(b):** Arrays of Confined
  Added a necessary restriction.

**(b) Between OTJLD 1.1 and OTJLD 1.2**

- **§ 1.2.2.(h):** Externalized creation
  Added alternative syntax using value parameter and changed title.

- **§ 1.2.5.(f):** Imports in role files
  Added a missing rule defining the effect of imports in role files.

- **§ 1.3.1.(c):** @Override annotation for roles
  The regular @Override annotation (Java ≥ 5) has been extended to apply to role classes, too.

- **§ 1.3.1.(k):** Covariant return types
  Necessary constraint for covariant return types in the presence of both implicit and explicit inheritance.

- **§ 2.1.2.(c):** Binding to final base class
  It has been added that binding to a final base class is now considered as decapsulation, too.
B.2. CHANGES BETWEEN VERSIONS OTJLD v1.3

• § 2.2.(f) : Ambiguous lowering
  A diagnostic has been added to detect situations where lowering might be intended but fails because the declared type is `java.lang.Object`, which makes a potential lowering translation unnecessary and thus ambiguous.

• § 2.3.2.(e) : Generic declared lifting
  Support passing unrelated base types into the same method with declared lifting.

• § 2.6.(g) : Decapsulation via base reference
  Extended applicability of decapsulation to two more positions.

• § 4.3.(f) : Base super call
  Support base calls directly to the super version of the bound base method, thus bypassing both the exact bound base method and also any further callins relating to this base method or its super version.

• § 5.4.(b) : Side-effects in guard predicates
  Migrate previous note about a future feature to a regular paragraph.

• § 5.4.(c) : Exceptions in guard predicates
  Clarify the effect of exceptions thrown from a guard predicate.

• § 6.2.(d) : LiftingVetoException
  Added documentation for the mostly internal `LiftingVetoException` and how it could actually be used in client code.

• § 6.2.(e) : Role migration
  Added two interfaces to add migration capabilities to a role class.

(c) Between OTJLD 1.2 and OTJLD 1.3

• § 2.1.1 : Binding roles to base interfaces
  The implementation limitation mentioned in § 2.1.1 has been mostly removed.

• § 2.3.1.(d) : Fine-tuning role instantiation
  An annotation has been defined for modifying the semantics of lifting in order to improve performance. Also a new section has been added as § 6.3 to summarize the annotation types defined in this document.

• § 2.3.5 : Consequences of lifting problems
  After § 2.3.4 has clarified that `LiftingFailedException` is indeed a checked exception, a subsection has been added defining the consequences of this exception in various program situations.

• § 3.1.(k) : Callout to generic method
  Added a rule on how a callout binding may refer to a generic base method.

• § 4.1.(b) : Callin binding in "unliftable" role
  Callin bindings can now be defined even in "unliftable" roles.

• § 4.1.(h) : Binding to team methods
  before and after callin bindings can now bind to methods of an enclosing class, too.
• § 4.8.(d): **Order when merging precedence declarations**
  Clarified how several precedence declarations are merged, which was underspecified, because the C3 algorithm needs ordered inputs, but this order was not specified.

• § 4.10.(e): **Propagating type parameters in callin bindings**
  In addition to capturing covariant return types, a callin binding may also declared type parameters in order to propagate genericity from its base method to the role method.

• § 5.3.(d): **Configuring implicit activation**
  Mechanisms have been added for configuring implicit team activation. The default has been changed to not apply implicit activation. A corresponding note has also been added to § 5.3.

• § 9.2.1.(a): **Instance constrained type parameter**
  Type anchors can now be applied to type parameters, too, thus expressing a new kind of constraint on the type parameter.